#### Trails of triples in Steiner triple systems

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Steiner triple systems and block colourings

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An STS(9) admitting a colouring of type (3, 3, 2, 2, 1, 1)

**Theorem** [Kirkman (1847)] An STS(v) exists if and only if  $v \ge 1$  and  $v \equiv 1, 3 \pmod{6}$ .

Partial Steiner triple systems and block colourings

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A PSTS(8)

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A PSTS(8) admitting a colouring of type (2, 2, 1, 1, 1)

For a PSTS(v) which admits a colouring of type  $(c_1, c_2, \ldots, c_t)$  to exist we must have

- (i)  $c_i \leq \lfloor \frac{v}{3} \rfloor$  for  $i = 1, 2, \ldots, t$ ; and
- (ii)  $c_1 + c_2 + \cdots + c_t \le \mu(v)$ , where  $\mu(v)$  is the maximum number of triples in a PSTS(v).

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- ► A Kirkman triple system is an STS(6t + 3) admitting a colouring of type (2t + 1, 2t + 1,..., 2t + 1).
- ► A nearly Kirkman triple system is a maximum PSTS(6t) admitting a colouring of type (2t, 2t, ..., 2t).
- A Hanani triple system is an STS(6t + 1) admitting a colouring of type (2t, 2t, ..., 2t, t).
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- The conjecture is also related to 3-frames and to many other block colouring problems.
- A strong Kirkman signal set SKSS(v, m) is a maximum PSTS(v) admitting a colouring of type (m, m, ..., m, r), where 1 ≤ r ≤ m.

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The realisability of many colour types follows immediately from the realisability of others.

For example, any STS(15) admitting a colouring of type (5, 5, 5, 5, 5, 5, 5, 5) must also admit a colouring of type (5, 5, 5, 5, 5, 5, 4, 3, 2, 1).

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But, for any large v, there are still vast numbers of feasible colour types which are not implied in this way.

For example, for v = 15, (4, 4, 4, 4, 4, 4, 4, 3) is not implied by (5, 5, 5, 5, 5, 5, 5) (or any other colour type).

If we can order the triples of a PSTS in such a way that any *m* consecutive triples are vertex disjoint, then the PSTS must admit all colourings of type  $(c_1, c_2, \ldots, c_t)$  where  $c_i \leq m$  for  $i = 1, 2, \ldots, t$ .

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Note that *m* can be at most  $\lfloor \frac{v}{3} \rfloor$ .

**Theorem** [Cohen, Colbourn (2000)] For each  $v \ge 1$  such that  $v \equiv 1, 3 \pmod{6}$ , there is an STS(v) admitting an  $\lfloor \frac{v+6}{9} \rfloor$ -pessimal ordering.

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**Corollary** For each sufficiently large v, each v-feasible colour type  $(c_1, c_2, \ldots, c_t)$  with  $c_i \leq m$  for  $i = 1, 2, \ldots, t$  is realisable.

# Proof sketch

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**Lemma** Let *n* be an odd integer. There exists an decomposition of the complete tripartite graph  $K_{n,n,n}$  into triples which admits an (n-2)-pessimal ordering.

# Proof of lemma



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# Example: an STS(73)

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### General case

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## Future directions

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- Can we improve our result to  $m = \frac{v}{3} O(1)$ ?
- Can we make this construction better by making it recursive?
- Prove the colouring conjecture.
- Latin square equivalents of these problems.
- For sufficiently large v, is there a maximum PSTS(v) that admits all v-feasible colourings?

# That's all.