

# Improved bounds on book crossing numbers of complete bipartite graphs via semidefinite programming

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# Outline

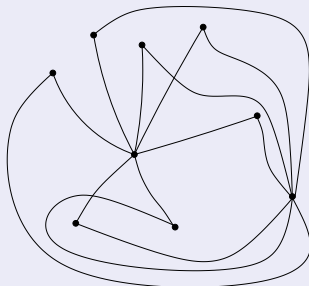
- The  $k$ -page crossing number of a complete bipartite graphs.
- Computing  $k$ -page crossing numbers by solving related  $\text{max-}k\text{-cut}$  problems.
- Semidefinite programming bounds and their implications.

# Crossing number of a graph

## Definition

The **crossing number**  $\text{cr}(G)$  of a graph  $G = (V, E)$  is the minimum number of edge crossings that can be achieved in a drawing of  $G$  in the plane.

## Example: the complete bipartite graph

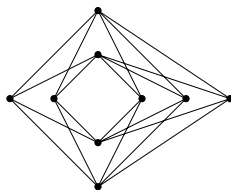


A drawing of  $K_{2,7}$  with 2 edge crossings. Not optimal, since  $\text{cr}(K_{2,7}) = 0$ .

# The Zarankiewicz conjecture

$K_{m,n}$  can be drawn in the plane with at most  $Z(m, n)$  edges crossing, where

$$Z(m, n) = \left\lfloor \frac{m-1}{2} \right\rfloor \left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n}{2} \right\rfloor.$$



A drawing of  $K_{4,5}$  with  $Z(4, 5) = 8$  crossings.

## Zarankiewicz conjecture (1954)

$$\text{cr}(K_{m,n}) \stackrel{?}{=} Z(m, n).$$

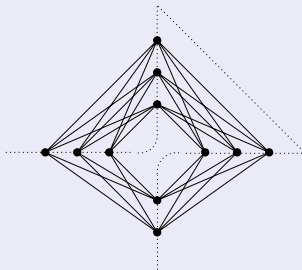
Known to be true for  $\min\{m, n\} \leq 6$  (Kleitman, 1970), and some special cases.

# $k$ -page crossing number of a graph

## Definition

In a  **$k$ -page (book) drawing** of  $G = (V, E)$  all vertices  $V$  must be drawn on a straight line (the spine of a book), and each edge in one of  $k$  half-planes incident to this line (the book pages). The  **$k$ -page crossing number  $\nu_k(G)$**  corresponds to  $k$ -page drawings of  $G$ .

Example: the complete bipartite graph  $K_{5,6}$



"Straighten the dotted line" to get a two-page drawing.

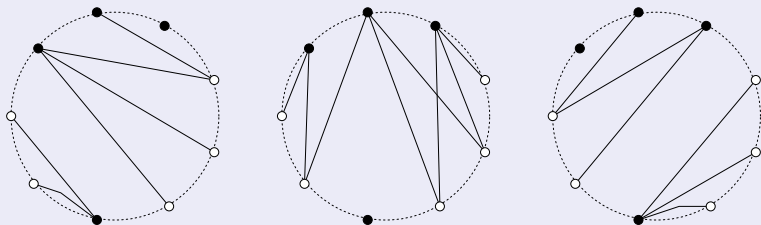
# Equivalent $k$ -page drawings

## Circular drawings

In a **circular drawing**, the vertices are drawn on a circle, and all edges inside the circle.

A  $k$ -page drawing is equivalent to  $k$  circular drawings.

Example: A 3-page drawing of  $K_{4,5}$  with 1 crossing.



# Applications and complexity

- Crossing numbers are of interest for graph visualization, VLSI design, quantum dot cellular automata, ...
- It is **NP-hard** to compute  $cr(G)$  or  $\nu_2(G)$  [Garey-Johnson (1982), Masuda et al. (1987)];
- The ( $k$ -page) crossing numbers of  $K_n$  and  $K_{n,m}$  are not known in general, ...
- Crossing number of  $K_{n,m}$  known as **Turán brickyard problem** — posed by Paul Turán in the 1940's.

Erdős and Guy (1973):

"Almost all questions that one can ask about crossing numbers remain unsolved."

Anno 2012 the situation has not changed much ... but there is some recent progress on **lower bounds** ...

# Some known results

## One page crossing number:

### Theorem (Riskin (2003))

If  $m|n$  then  $\nu_1(K_{m,n}) = \frac{1}{12}n(m-1)(2mn-3m-n)$ , and this minimum value is attained when the  $m$  vertices are distributed evenly amongst the  $n$  vertices.

## Two page crossing number:

### Theorem (De Klerk, Pasechnik, Schrijver (2007))

$$1 \geq \lim_{n \rightarrow \infty} \frac{\nu_2(K_{m,n})}{Z(m,n)} \geq \lim_{n \rightarrow \infty} \frac{ct(K_{m,n})}{Z(m,n)} \geq 0.8594 \text{ if } m \geq 9.$$

### Theorem (De Klerk and Pasechnik (2011))

$$\lim_{n \rightarrow \infty} \frac{\nu_2(K_{m,n})}{Z(m,n)} = 1 \text{ if } m \in \{7, 8\}.$$

These results use **SDP lower bounds**.



# Some known results (ctd.)

## **$k$ -page crossing number:**

Theorem (Shahrokhi et al. 1996 (lower bound); De Klerk, Pasechnik, Salazar 2012 (upper bound))

$$\frac{1}{3(3\lceil \frac{k}{2} \rceil - 1)^2} \leq \lim_{m,n \rightarrow \infty} \frac{\nu_k(K_{m,n})}{\binom{m}{2} \binom{n}{2}} \leq \frac{1}{k^2}.$$

# New results (this talk)

Theorem (De Klerk, Pasechnik, Salazar (2012))

Let  $k \in \{2, 3, 4, 5, 6\}$ , and let  $n$  be any positive integer. Define  $\ell := \left\lfloor \frac{(k+1)^2}{4} \right\rfloor$  and  $q := n \bmod \left\lfloor \frac{(k+1)^2}{4} \right\rfloor$ . Then

$$\nu_k(K_{k+1,n}) = q \cdot \binom{\frac{n-q}{\ell} + 1}{2} + (\ell - q) \cdot \binom{\frac{n-q}{\ell}}{2}.$$

Also, asymptotically ( $k \rightarrow \infty$ ),

$$\lim_{k \rightarrow \infty} \left( \lim_{n \rightarrow \infty} \frac{\nu_k(K_{k+1,n})}{2n^2/k^2} \right) = 1.$$

# New result: outline of the proofs

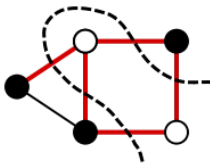
- We sketch the proof of the first result (lower bound only).
- One may compute  $\nu_k(K_{m,n})$  by solving some **maximum  $k$ -cut problems** (Buchheim and Zheng (2007));
- We first show, for  $2 \leq k \leq 6$ , that  $\nu_k(K_{k+1,s}) > 0$  for some  $s = s(k)$ .
- This is done by computing **semidefinite programming bounds**.
- Then we use a result by Turán to obtain lower bounds on  $\nu_k(K_{k+1,n})$  **for general  $n$** .

# The maximum $k$ -cut problem for graphs

## Definition

For  $G = (V, E)$  and a set of  $k$  colors, color  $V$  such that the number of edges with differently colored end points is a maximum. This maximum is denoted by  $\max\text{-}k\text{-cut}(G)$ .

**Example** ( $k = 2$ ):



A maximum 2-cut example with  $\max\text{-}k\text{-cut}(G) = 5$ .

# $\nu_k(K_{m,n})$ and the maximum $k$ -cut problem

$\nu_k(K_{m,n})$  may be obtained by solving some **maximum  $k$ -cut problems** ... one for each possible one-page (circle) drawing  $D$  of  $K_{m,n}$ .

Given  $D$ , the max- $k$ -problem is solved for the graph  $G_D(K_{m,n}) = (V_D, E_D)$  where:

- $V_D$  is the edge set of  $K_{n,m}$ ;
- Two vertices in  $V_D$  are adjacent **if the corresponding edges cross** in  $D$ .

**Lemma** (cf. Buchheim and Zheng (2007))

One has

$$\nu_k(K_{m,n}) = \min_D (|E_D| - \max\text{-}k\text{-cut}(G_D(K_{n,m}))).$$

**Proof:** Given a  $k$ -page (circle) drawing of  $K_{m,n}$ , assign the edges drawn on page  $i$  the color  $i$  ( $1 \leq i \leq k$ ).

# The number of one page drawings of $K_{m,n}$

**Q:** How many distinct one-page drawings are there of  $K_{m,n}$ ?

**A:** as many as there are **orbits** of the dihedral group  $D_{m+n}$  acting on the set of one-page drawings.

Thus we may use the **Burnside 'orbit counting' lemma**.

## Lemma (Orbit counting lemma (Frobenius?))

*Let a finite group  $\mathcal{G}$  act on a finite set  $\Omega$ . Denote by  $\Omega^g$ , for  $g \in \mathcal{G}$ , the set of elements of  $\Omega$  fixed by  $g$ . Then the number  $N$  of orbits of  $\mathcal{G}$  on  $\Omega$  equals*

$$N = \frac{1}{|\mathcal{G}|} \sum_{g \in \mathcal{G}} |\Omega^g|.$$

# The number of one page drawings of $K_{m,n}$

Applying the orbit counting lemma we obtain:

## Lemma

Let  $d = \gcd(m, n)$ . The number of distinct circular drawings of  $K_{m,n}$  equals:

$$\frac{1}{2(m+n)} \begin{cases} \frac{m+n}{2} \left( \binom{\frac{m+n}{2}}{n/2} + \binom{\frac{m+n-2}{2}}{m/2} + \binom{\frac{m+n-2}{2}}{n/2} \right) + \sum_{k=0}^{d-1} \binom{\frac{m+n}{o(k)}}{\frac{m}{o(k)}} & (m, n \text{ even}) \\ (m+n) \binom{\frac{m+n-1}{2}}{n/2} + \sum_{k=0}^{d-1} \binom{\frac{m+n}{o(k)}}{\frac{m}{o(k)}} & (m \text{ odd}, n \text{ even}) \\ (m+n) \binom{\frac{m+n-2}{2}}{(m-1)/2} + \sum_{k=0}^{d-1} \binom{\frac{m+n}{o(k)}}{\frac{m}{o(k)}} & (m, n \text{ odd}) \end{cases}$$

where  $o(k)$  is the order of the subgroup generated by  $k$  in the additive group of integers mod  $d$ .

**Example:** there are **1980** distinct one-page drawings of  $K_{7,13}$ .

# Bounds from the $\vartheta$ -function

## Lemma (Lovász (1979))

Given a graph  $G = (V, E)$  and the value

$$\vartheta(G) := \max_{X \succeq 0} \left\{ \sum_{i,j \in V} X_{ij} \mid X_{ij} = 0 \text{ if } (i,j) \in E, \text{ trace}(X) = 1, X \in \mathbb{R}^{V \times V} \right\},$$

one has

$$\omega(\bar{G}) \leq \vartheta(G) \leq \chi(\bar{G}),$$

where  $\omega(\bar{G})$  and  $\chi(\bar{G})$  are the clique and chromatic numbers of the complement  $\bar{G}$  of  $G$ , respectively.

## Corollary

If  $\vartheta(\overline{G_D(K_{m,n})}) > k$  for all one-page drawings  $D$  of  $K_{m,n}$ , then  $\nu_k(K_{m,n}) > 0$ .



# Bounds from the $\vartheta$ -function (ctd)

We had:

## Corollary

*If  $\vartheta(\overline{G_D(K_{m,n})}) > k$  for all one-page drawings  $D$  of  $K_{m,n}$ , then  $\nu_k(K_{m,n}) > 0$ .*

By computing the  $\vartheta$ -function for all distinct one-page drawings using DSDP (Benson-Ye) we obtained:

## Theorem

*For each  $k \in \{2, 3, 4, 5, 6\}$ ,*

$$\nu_k(K_{k+1, \lfloor (k+1)^2/4 \rfloor + 1}) > 0.$$

(For a few drawings, where  $\vartheta(\overline{G_D(K_{m,n})}) = k$ , we had to compute the chromatic number **exactly** using XPRESS-MP.)

We can use these results to obtain lower bounds on  $\nu_k(K_{k+1,n})$  ( $k \in \{2, 3, 4, 5, 6\}$ ) for **general**  $n$ , by using **Turán's theorem**.

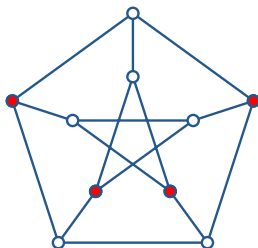
# Turán's theorem

## Theorem (Turán (1941))

If  $G = (V, E)$  has coclique number  $\alpha(G) \leq s$  for some  $s$ , then

$$|E| \geq \lceil (1/2)|V|^2(1/s - 1/|V|) \rceil.$$

**Example:** Petersen graph:  $|V| = 10$ ,  $|E| = 15$ ,  $\alpha(G) = 4$ .



$$|E| \geq \lceil (1/2)10^2(1/4 - 1/10) \rceil = 8.$$

# Applying Turán's theorem

- 1 Assume  $\nu_k(K_{k+1,s}) > 0$  for some  $k, s$ .
- 2 Let  $D$  be an **optimal  $k$ -page drawing of  $K_{k+1,n}$**  for some  $n > s$ .
- 3 Now construct an auxiliary  $G_D = (V_D, E_D)$ :  $V_D$  is the vertices from the  $n$  co-clique of  $K_{k+1,n}$ . Two vertices in  $V$  are adjacent if **two edges incident to them cross** in  $D$ .
- 4 **NB:**  $\nu_k(K_{k+1,n}) \geq |E_D|$  and  $\alpha(G_D) \leq s$ .
- 5 Now apply **Turán's theorem** to obtain:

$$\nu_k(K_{k+1,n}) \geq (1/2)n^2(4/(k+1)^2 - 1/n) \quad (2 \leq k \leq 6, n \geq (k+1)^2/4).$$

# Conclusion and summary

- We demonstrated **improved lower bounds** on the  $k$ -page crossing numbers of some complete bipartite graphs.
- The proofs for small  $k$  were **computer-assisted**, using semidefinite programming (SDP) relaxations.
- Preprints available at **arXiv**.