# Reconstructing phylogenetic level-1 networks from nondense binet and trinet sets

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#### **Supertrees**



#### **Supernetworks**



## **Trinets and Subnets**

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#### The *subnet* N|X' is obtained by

- 1. taking all directed paths in N from LSA(X') to a leaf in X';
- 2. suppressing indegree-1 outdegree-1 vertices and parallel arcs.



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A *trinet* is a subnet with 3 leaves. A *binet* is a subnet with 2 leaves.



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**Definition.** (assuming binary reticulations)

- *level*-*k*: at most *k* reticulations per biconnected component;
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**Definition.** (assuming binary reticulations)

- *level*-*k*: at most *k* reticulations per biconnected component;
- *tree-child*: each non-leaf vertex has a child that is not a reticulation.

**Theorem.** (Huber, vI & Moulton) Binary level-1, level-2 and tree-child networks are all encoded by their trinets.

#### **Reconstruction Algorithms**

Given any set of triplets, we can construct a tree displaying them, if one exists, in polynomial time. (Aho, Sagiv, Szymanski, Ullman, 1981)



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Given any set of *binets* we can construct a *level-1 network* displaying them, if one exists, in polynomial time. (Huber, vI, Moulton, Scornavacca, Wu, 2015)



**Step 1:** can the network have a root that is not in a cycle?

a, b, c, d, e f, g

N



Graph  $\mathscr{R}$  connects taxa that have to be on the same side of the root.

*Step 2:* if the root is in a cycle, which taxa are *low* (below the cycle), and which ones are *high* (to the side of the cycle)?



Graph  $\mathcal{K}$  connects taxa that have to be *at the same height*.

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Find a proper subset of the vertices with no incoming arcs. Make those taxa "high".

Digraph  $\Omega$  has arcs indicating which taxa have to be *above* other taxa.



**Step 3:** recursively find the sidenetworks



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... but this problem is polynomial-time solvable for subnets with only size-3 cycles (the blue ones) ...

... and there is an  $O(3^n \text{poly}(n))$  time algorithm for the general case, with *n* taxa.

- 1. Is the root in a cycle?
- 2. Which taxa are high/low?
- 3. Which taxa are on the left/right?
- 4. Partition each side into sidenetworks.
- 5. Recursively find all sidenetworks.



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N

a, b, d, e

*Step 3*: which taxa are on the left, which taxa on the right?



h

Graph  $\mathcal{M}$  connects taxa that have to be **on the same side**.

a **'** 

*Step 3*: which taxa are on the left, which taxa on the right?



Graph *W* connects components that have to be *on different sides*.

*Step 3*: which taxa are on the left, which taxa on the right?











*Step 5:* recursively find sidenetworks.



- Constructing a level-1 supernetwork from trinets is NP-hard.
- Our exponential-time algorithm heavily exploits the structure of the problem.
- Some special cases can be solved in polynomial-time.
  - binets;
  - subnets with only size-3 cycles.



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