

# Zero-sum Revision Games

Fabien Gensbittel<sup>1</sup>, Stefano Lovo<sup>2</sup>, Jerome Renault<sup>1</sup> and  
Tristan Tomala<sup>2</sup>

<sup>1</sup>TSE

<sup>2</sup>HEC, Paris

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- Players have to prepare their actions in a pre-play phase preceding the payoff-relevant play in a one shot game,
- During the pre-play phase:
  - prepared actions are commonly observed.
  - Prepared actions can be change only at the bell of a Poisson clock.
- Only the last prepared action profile matters for the payoff.

## Some examples

- Preopening in the stock market (Nasdaq, Euronext, Toronto SE, daily from 7a.m. to 9 a.m.)
- Interaction through internet servers (e-bay auctions).
- Preparatory meetings to negotiate the terms of a treaty.
- Armies deploying their troops on the ground
- ...

# Component game: Zero-sum game

- 2 players.
- $X_i$ : player  $i$ 's finite set of actions.
- $U : X_1 \times X_2 \rightarrow \mathbb{R}$ : player 1's payoff matrix (generic).
- 

$$BR_1^U(x) := \arg \max_{y_1} U(y_1, x_2) ; BR_2^U(x) := \arg \min_{y_2} U(x_1, y_2)$$

- Stackelberg payoff where 1 plays first:

$$S_1 = \max_{x_1 \in X_1} \min_{x_2 \in X_2} U(x_1, x_2)$$

- Stackelberg payoff where 2 plays first

$$S_2 = \min_{x_2 \in X_2} \max_{x_1 \in X_1} U(x_1, x_2)$$

- Value of the game:  $V$

$$S_1 \leq V \leq S_2$$

with equality if pure Nash.

# Revision game

- At  $t = 0$ , starting prepared action  $x(0) \in X$  exogenous.
- Between time 0 and  $T$ , Poisson arrivals of revision times independent for each player (Asynchronous moves).
- Each player can change his prepared actions only at his revision times.
- At  $T$  players get their only payoff and this results from players playing, in the component game, their last prepared actions.

# Revision game as a stochastic game

- Finite time horizon  $[0, T]$ .
- Game  $\Gamma_{[\tau, T]}(x)$  with  $\tau < T$  and  $x \in X$ .
- Time  $\eta$  is drawn from an exponential distribution with parameter  $\lambda$ .
- If  $\eta + \tau > T$ , then the game is over and players' payoff is  $\{U(x), -U(x)\}$
- if  $\eta + \tau < T$ , then
  - With Pr  $q \in (0, 1)$  player 1 chooses an action  $y_1 \in X_1$  and the game  $\Gamma_{[\tau+\eta, T]}(y_1, x_2)$  starts.
  - With Pr  $1 - q$  player 2 chooses an action  $y_2 \in X_2$  and the game  $\Gamma_{[\tau+\eta, T]}(x_1, y_2)$  starts.
- Initial game:  $\Gamma_{[0, T]}(x(0))$

- **Non-zero sum revision game:** Kamda and Kandori 2009; Lovo and Tomala (2015).

Calcagno, Kamada, Lovo and Sugaya (2014): In  $2 \times 2$  conflicting interest games (generic Battle of the sexes),

- the revision game equilibrium is unique;
  - the slow players has an advantage over the fast players;
  - revision game equilibrium payoff = component game Nash equilibrium payoff;
  - All action occurs at the beginning of the revision game.
- **Cheap talk games:** Farrell (1987), Rabin (1994), Aumann and Hart (2003), ...
  - **Switching cost games:** Lipman and Wang (2000) and Caruana and Einav (2008), ...

## What are we after?

- Under what conditions does a player prefer to play the revision game rather than the straight zero-sum game?
- Characterization of equilibrium payoff.
- Characterization of equilibrium behavior.

- 1 Preliminaries
- 2 General results
- 3  $2 \times 2$  equilibrium chracterization



# Preliminaries

## Notation, histories and strategies

- Set of states: player who can revise and the resulting new profile of action

$$K = \{1, 2\} \times X$$

- History of past revision time and chosen actions

$$h_n = \{x, \tau_1, k_1, \dots, \tau_n, k_n\} \in X \times ([0, T] \times K)^n$$

- Strategy: mapping histories and revision times into a (mixed) action

$$\sigma_i : \cup_{n \geq 0} (H_n \times [0, T]) \rightarrow \Delta X_i$$

- A Markov strategy is a measurable mapping

$$\sigma_i : X \times [0, T] \rightarrow \Delta X_i$$

- Expected payoff given  $\sigma$ :

$$u_\sigma(T, x) := \mathbb{E}_\sigma[U(x(T)) | x(0) = x]$$

where  $x(T)$  is the last prepared action profile at time  $T$ .

## Theorem

*(Lovo and Tomala (2015)) The revision game has a Markov perfect equilibrium. With*

- *$t$  is the remaining time.*
- *$u(x, t)$ : equilibrium payoff of the game of length  $t$  with starting action profile  $x$ .*
- *$u(t) := \{u(t, x)\}_{x \in X}$*
- *$\sigma_i(t, x) \in BR_i^{u(t)}(x_{-i})$*

# Preliminaries

## Dynamic programming equation

Remark that  $u(t, x)$  is Lipschitz.

Let

$$u^+(t, x) := \max_{y_1 \in X_1} u(t, y_1, x_2) \quad ; \quad u^-(t, x) := \min_{y_2 \in X_2} u(t, x_1, y_2);$$
$$\lambda_1 := \lambda q \quad ; \quad \lambda_2 := \lambda(1 - q)$$

Then

$$u(t, x) = U(x)e^{-\lambda t} + \int_{s=0}^t e^{-\lambda(t-s)} (\lambda_1 u^+(s, x) + \lambda_2 u^-(s, x)) ds,$$
$$\frac{\partial u(t, x)}{\partial t} = \lambda_1(u^+(t, x) - u(t, x)) + \lambda_2(u^-(t, x) - u(t, x)),$$
$$u(0, x) = U(x).$$

# General results

## General equilibrium properties

### Proposition

- ① *The revision game has a Markov perfect equilibrium in pure strategy.*
- ② *The equilibrium payoff  $u(t, x)$  is Lipschitz in  $t$ ,  $U$  and is continuous in  $(q, \lambda) \in (0, 1) \times (0, \infty)$ .*
- ③ *The equilibrium payoff  $u(t, x)$  is non-decreasing in  $q$ .*

# General results

Pure strategies: sketch of the proof

Take a MPE and suppose that for some time  $t$ ,  $\sigma_i(t, x)$  is not pure. For this  $t$  replace  $\sigma_i(t, x)$  by a pure action in  $\sigma'_i(t, x) \in BR_i^{u(t)}(x)$ . Observe that  $u^+$  and  $u^-$  do not change with  $\sigma$  or  $\sigma'$ . Hence

$$u(t, x) = u(x)e^{-\lambda t} + \int_{s=0}^t e^{-\lambda(t-s)} (\lambda_1 u^+(s, x) + \lambda_2 u^-(s, x)) ds$$

is the same under  $\sigma$  and  $\sigma'$ .  
Zero sum structure is crucial.

# General results

Continuity of  $u(t, x)$ : sketch of the proof

- 1-Lipschitz in  $t$ :

$$|u(t, x) - u(t + \varepsilon, x)| \leq \|u(t)\| (1 - e^{-\lambda\varepsilon})$$

- 1-Lipschitz in  $U$ : Take  $U' \neq U$ , then

$$|u(t, x) - u'(t, x)| \leq \max_{y \in X} |U(y) - U'(y)|$$

- Continuous in  $\lambda$ , take  $\lambda' \neq \lambda$ , then

$$u(t, x)|_{\lambda} = u\left(\frac{\lambda'}{\lambda}t, x\right)\Big|_{\lambda'}$$

- Continuous in  $q$ : Payoff continuously depends on the distribution of revision time that is continuous in  $q$ .

# General results

Monotonicity of  $u(t, x)$  in  $q$ : sketch of the proof

Let  $q' < q$ .

$$\frac{\partial u(t, x)}{\partial t} = \lambda(qu^+(t, x) + (1 - q)u^-(t, x) - u(t, x))$$

If for some  $t$ ,  $u(t, x)|_{q'} = u(t, x)|_q$ , then

$$\left. \frac{\partial u(t, x)}{\partial t} \right|_{q'} \leq \left. \frac{\partial u(t, x)}{\partial t} \right|_q$$

implying  $u(\tau, x)|_q \leq u(\tau, x)|_{q'}$ , for some  $\varepsilon > 0$ , and all  $\tau \in (t, t + \varepsilon)$ .

but

$$u(0, x)|_{q'} = u(0, x)|_q = U(x)$$

so it can never be that

$$u(\tau, x)|_{q'} > u(\tau, x)|_q$$

# General results

## Revision game value

Consider a revision game where the starting action profile is  $x$  and let

$$\underline{R}(x) := \liminf_{t \rightarrow \infty} u(t, x) \quad \text{and} \quad \overline{R}(x) := \limsup_{t \rightarrow \infty} u(t, x)$$

If  $\underline{R}(x) = \overline{R}(x) = R(x)$  then we say that the *revision game value* is  $R(x)$



# General results

## Properties of the revision game value

### Proposition

- ① *Irrelevance of revision when  $V$  is achieved with pure strategies:*

$$S_1 \leq \underline{R}(x) \leq \overline{R}(x) \leq S_2$$

- ② *Ergodicity:*

$$\underline{R}(x) = \overline{R}(x) = R, \forall x \in X$$

*for some constant  $C$*

- ③  *$R$  is 1-Lipschitz in  $U$ , and continuous in  $(q, \lambda) \in (0, 1) \times (0, \infty)$ .*

④

$$\lim_{q \rightarrow 1} = S_2 \text{ and } \lim_{q \rightarrow 0} = S_1$$

# 2 × 2 games

Let  $X_1 = X_2 = \{\alpha, \beta\}$ .

	$\alpha$	$\beta$
$\alpha$	$U(\alpha, \alpha)$	$U(\alpha, \beta)$
$\beta$	$U(\beta, \alpha)$	$U(\beta, \beta)$

Then the component game where  $U$  is generic, implying  $x \neq x' \Rightarrow U(x) \neq U(x')$ .

# 2 × 2 games

Possible scenarios for  $u(t)$

- **Scenario DD:** Each player has a dominant action. For example:

A 2x2 game matrix with arrows indicating dominant strategies. The matrix is:

	←		
↑	0	2	↑
	-1	1	
	←		

- **Scenario DN:** One player has a dominant action whereas the other player does not. For example:

A 2x2 game matrix with arrows indicating dominant strategies. The matrix is:

	←		
↑	0	2	↑
	-1	-5	
	→		

- **Scenario NN:** No pure Nash Eq. For example:

A 2x2 game matrix with arrows indicating no pure Nash equilibrium. The matrix is:

	←		
↓	0	2	↑
	3	-5	
	→		

### Proposition

*Suppose  $U$  is in scenario DD, and let  $\hat{x}_i$  be player  $i$ 's dominant action in the component game. Then for all  $t$ ,*

①  *$u(t)$  is in scenario DD*

②

$$\sigma_i(t, x) = \hat{x}_i, \forall x \in X$$

③

$$R = V$$

# 2 × 2 games

Scenario DD: Sketch of the proof

## Intuition:

- By continuity with respect to  $t$  each player prepares his dominant action when  $t$  is close to 0.
- If for all  $\tau > t$  the other player uses a fixed action no matter what you do, then you strictly prefer preparing your dominant action at  $t$ .

## Algebraic:

Solve the ODE and verify that  $BR_i^{u(t)}(x)$  does not depend on  $t$ .

# 2 × 2 games

## Scenario DN

### Proposition

*Suppose  $U$  is in scenario DN, and  $\hat{x}_1$  is player 1's dominant action in the component game. Then there is  $t^*$  finite such that*

① For  $t < t^*$ :

- $u(t)$  is in scenario DN
- $\sigma_1(t, x) = \hat{x}_1, \forall x$
- $\sigma_2(t, x) = BR_2^U(x_1)$

② For  $t \geq t^*$ :

- $u(t)$  is in scenario DD
- $\sigma_1(t, x) = \hat{x}_1, \forall x$
- $\sigma_2(t, x) = BR_2^U(\hat{x}_1), \forall x$

③

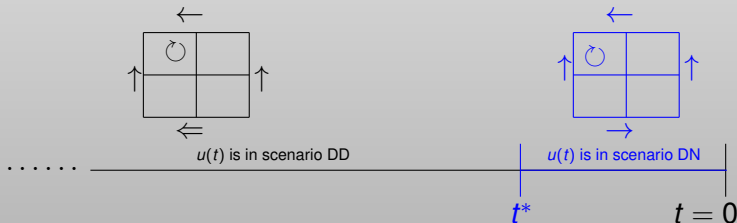
$$R = V$$

# 2 × 2 games

## Scenario DN

	$\alpha$	$\beta$
$\alpha$	0	2
$\beta$	-1	-5

- 1  $t > t^*$ : At the beginning of the revision phase players prepare the action forming the component game pure Nash equilibrium.
- 2  $t < t^*$ : Once reached these actions they do not move.



### Proposition

If  $U$  is in NN scenario, then there are  $0 < t^{**} < t^*$ ,  $i^*$  and  $x_{i^*}^*$  such that:

- ① For  $t < t^{**}$ :
  - $u(t)$  is in scenario NN
  - $\sigma_i(t, x) = BR_i^U(x_{-i})$
- ② For  $t^{**} < t < t^*$ :
  - $u(t)$  is in scenario DN
- ③ For  $t \geq t^*$ :
  - $u(t)$  is in scenario DD
- ④ Generically

$$R = u(t^{**}, x^*) \neq V$$



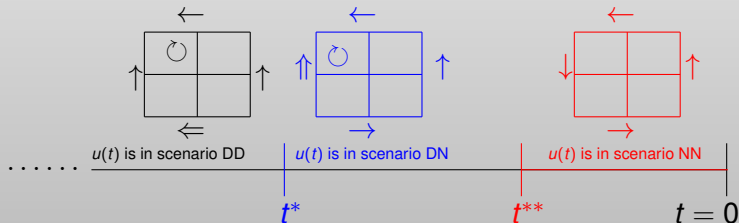
# 2 × 2 games

## Scenario NN

surplice and wrestle

0	0.2
0.3	-0.5

- 1  $t > t^{**}$ : At the beginning of the revision phase players prepare a “surplice” action, that they keep until  $t^*$
- 2  $t < t^*$ : starting to  $t^*$  players actions cycle



# 2 × 2 games

Scenario NN: Payoffs normalization

If  $U$  is in NN scenario, then without loss of generality we have

	$\alpha$	$\beta$
$\alpha$	0	$b$
$\beta$	$c$	$b + c - 1$

with

$$0 < b, c, < 1$$

# 2 × 2 games

Scenario NN: Sur-place action

Let  $\sigma := |\lambda_1^2 - 6\lambda_1\lambda_2 + \lambda_2^2|^{\frac{1}{2}}$  and let  $\hat{t}(A, B, \lambda_1, \lambda_2)$  be the smallest positive  $t$  such that

$$e^{-\frac{\lambda_1 + \lambda_2}{2}t} + A \cos\left(\frac{\sigma t}{2}\right) + \frac{(\lambda_2 - \lambda_1)A + 2\lambda_2 B}{\sigma} \sin\left(\frac{\sigma t}{2}\right) = 0 \quad (1)$$

Set  $\hat{t}(A, B, \lambda_1, \lambda_2)$  to infinity. Let

$$t_{\alpha, \alpha} = \hat{t}(2c - 1, 2b - 1, \lambda_1, \lambda_1)$$

$$t_{\alpha, \beta} = \hat{t}(2b - 1, 1 - 2c, \lambda_2, \lambda_1)$$

$$t_{\beta, \alpha} = \hat{t}(1 - 2b, 2c - 1, \lambda_2, \lambda_1)$$

$$t_{\beta, \beta} = \hat{t}(1 - 2c, 1 - 2b, \lambda_1, \lambda_1)$$

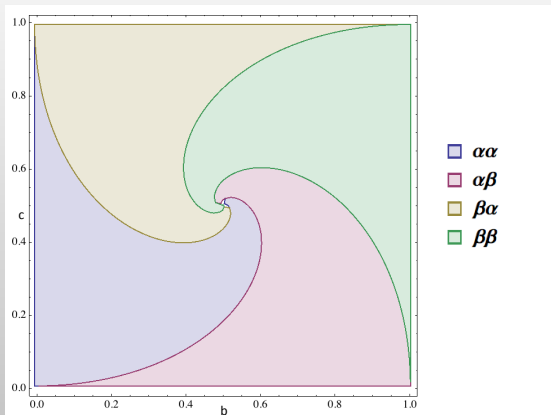
Then then

$$\hat{x} = \arg \min_{y \in \{(\alpha, \alpha), (\alpha, \beta), (\beta, \alpha), (\beta, \beta)\}} t_y$$

$$t_* = t_{\hat{x}}$$

# $2 \times 2$ games

Scenario NN: sur-place actions for  $q = 1/2$  and  $0 < b, c < 1$



	$\alpha$	$\beta$
$\alpha$	0	$b$
$\beta$	$c$	$b + c - 1$

# 2 × 2 games

Scenario NN:  $R$  and  $V$  for  $q = 1/2$  and  $0 < b, c < 1$

## Theorem

If  $0 < b, c < 1$  and  $q = 1/2$ , then

- The value of the game is  $V = bc$
- the revision game value is:

$$R = \frac{1}{4}(2c + 2b - 1) + \frac{1}{2}(c + b - 1)(b - c) \sin(2\mu) \\ + \frac{1}{4}(2b - 1)(2c - 1) \cos(2\mu),$$

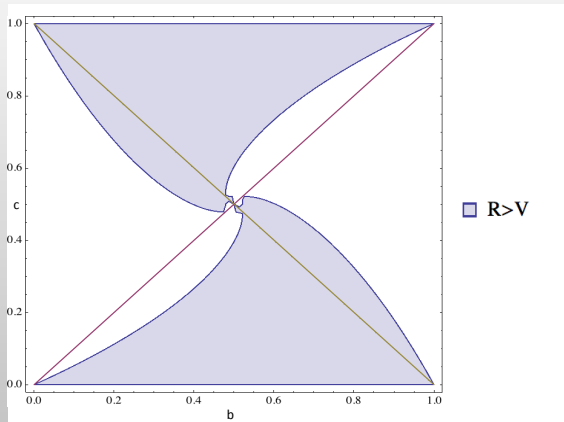
where  $\mu$  is the smallest  $t$  in  $\mathbb{R}_+$  satisfying:

$$e^{-t} = \max\{(1 - 2c) \cos(t) + (1 - 2b) \sin(t), (1 - 2b) \cos(t) - (1 - 2c) \sin(t),$$

$$-(1 - 2b) \cos(t) + (1 - 2c) \sin(t), -(1 - 2c) \cos(t) - (1 - 2b) \sin(t)\}.$$

# 2 × 2 games

Scenario NN:  $R$  and  $V$  for  $q = 1/2$  and  $0 < b, c < 1$

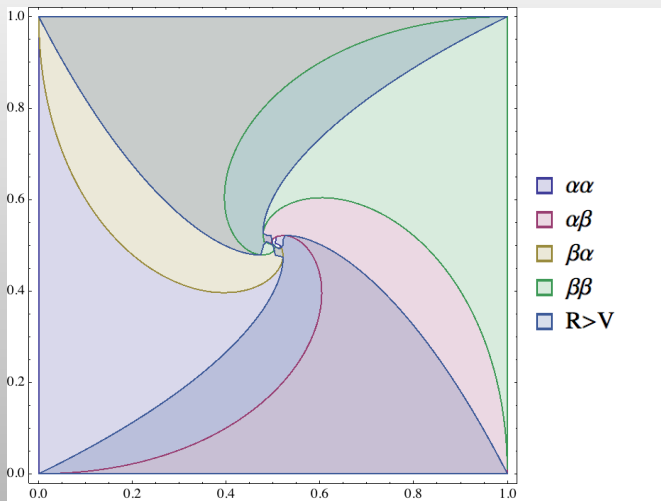


	$\alpha$	$\beta$
$\alpha$	0	$b$
$\beta$	$c$	$b + c - 1$

# $2 \times 2$ games

Scenario NN: Sur-place action,  $R$  and  $V$

$$q = 1/2$$



# Conclusion

- A zero-sum revision game always has a pure strategy equilibrium.
- When the component game Nash equilibrium is in pure, then players should be indifferent between playing the game with or without a (long) revision phase.
- When the component game Nash equilibrium is not pure, then
  - A player gain from being faster than the other player.
  - Generically the revision game value is different from the one-shot game value.
  - For  $2 \times 2$  games, the unique equilibrium consists in players waiting on a sur-place action profile until the the deadline is close and then wrestle.