

# Brown's lemma in reverse mathematics

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For instance, the 2-coloring

$$01001100011100001111\dots$$

has 1-piecewise syndetic homogeneous sets for both colors but no syndetic homogeneous sets.

The 2-coloring

$$010101010101\dots$$

has 2-syndetic homogeneous sets in both colors but no 1-piecewise syndetic homogeneous set.

The 2-coloring

$$0011001100110011\dots$$

has 3-syndetic homogeneous sets but no 2-piecewise syndetic homogeneous sets.

The classical proof of Brown's lemma is elementary and easily formalizable in second-order arithmetic. However, it is based on the following variant of König's lemma, which is easily seen to be equivalent to  $\text{ACA}_0$ .

### Lemma

*If  $S \subseteq r^{<\mathbb{N}}$  is infinite, then there exists  $g: \mathbb{N} \rightarrow r$  such that for all  $n$  there is  $\sigma \in S$  with  $g \upharpoonright n \subseteq \sigma$ .*

### Proposition

*$\text{ACA}_0$  is equivalent to  $\Sigma_1^0\text{-WKL}_0$ .*

## Theorem (F)

Over  $\text{RCA}_0$  (even  $\text{RCA}_0^*$ ), Brown's Lemma is equivalent to  $\text{I}\Sigma_2^0$ .

Piecewise syndetic sets are  $\Sigma_3$ -definable and closed under supersets. Brown's lemma does not actually assert the existence of a set. It is a  $\Pi_1^1$ -statement. A partition lemma of this form has no computational strength but only inductive strength.

We conjecture that a statement of the form

† Every finite coloring  $f: \mathbb{N} \rightarrow r$  has an  $f$ -homogeneous *large* set,

where *large* is a property about sets closed under supersets, is equivalent to

- $\text{I}\Sigma_n^0$  if *large* is  $\Sigma_{n+1}$ -definable,
- $\text{B}\Sigma_n^0$  if *large* is  $\Pi_n$ -definable.

## Evidence

Brown's lemma and the infinite pigeonhole principle ( $RT^1$ ) serve as an example for  $n = 2$ .

The following is equivalent to  $IS_2^0$  (Hirst):

- For every finite coloring  $f: \mathbb{N} \rightarrow r$  there is  $b$  such that for all  $x > a$  there is  $y > x$  such that  $f(x) = f(y)$ .

The following version of Folkman's theorem is equivalent to  $B\Sigma_2^0$ :

### Theorem

*Every  $f: \mathbb{N} \rightarrow r$  has an  $f$ -homogeneous set  $X$  such that for every  $k$  there exists a set  $F \subseteq X$  of size  $k$  with  $FS(F) \subseteq X$ .*

The proof uses the finite version of Folkman's theorem.

## Brown's Lemma vs van der Waerden's Theorem

In "Ramsey theory on the integers" page 279:

*Comparing this statement to the statements given in Theorem 2.5, we see that Brown's lemma is very reminiscent of van der Waerden's theorem. However, it is known that Brown's lemma neither implies, nor is implied by, van der Waerden's theorem*

I asked the authors what that means.

*By this we mean that assuming only BL, we cannot prove VDW (as you state you need VDW finite version to show infinite BL implies infinite VDW) and assuming only VDW, we cannot prove BL. So, even through infinite BL may be stronger than infinite VDW, the infinite BL does not by itself imply infinite VDW (or even finite VDW).*

Here VDW is the the finite one.



## Theorem (Van Der Waerden's Theorem)

*Every finite coloring  $f: \mathbb{N} \rightarrow r$  has an  $f$ -homogenous set with arbitrarily long arithmetic progressions.*

Let  $N = \{0, 1, \dots, N - 1\}$ .

## Theorem (Van Der Waerden's Theorem, finite)

*For all  $r, l$  there exists  $N$  such that every coloring  $C: N \rightarrow r$  has a  $C$ -homogeneous arithmetic progression of length  $l$ .*

Shelah provided primitive recursive upper bounds for the van der Waerden numbers as a byproduct of a new elementary proof of Hales-Jewett theorem that uses  $\Sigma_1$ -induction.

## Theorem (Folklore)

*The finite version of van der Waerden's theorem is provable in  $\text{RCA}_0$ .*

Shelah's proof of Hales-Jewett theorem avoids the use of double induction which results in Ackermannian upper bounds. Gowers in his celebrated work on Szemerédi's theorem obtained elementary recursive upper bounds for the van der Waerden's numbers. Gowers' bound is the following:

$$W(r, l) \leq 2^{2^{f(r, l)}}, \text{ where } f(r, l) = r^{2^{2^{l+9}}}.$$

However, the proof is far from elementary and so we can only conjecture that van der Waerden's theorem is provable in EFA (Elementary Function Arithmetic). On the other hand, the lower bounds for the van der Waerden's numbers are exponential, and hence van der Waerden's theorem is not provable in bounded arithmetic.

By the finite van der Waerden's theorem, every piecewise syndetic set contains arbitrarily long arithmetic progressions and this implication can be proved within  $\text{RCA}_0$ . Therefore it is not surprising that within  $\text{RCA}_0$  Brown's lemma implies van der Waerden's theorem.

We establish this by showing:

### Proposition

*Over  $\text{RCA}_0$  (even  $\text{RCA}_0^*$ ), van der Waerden's theorem is equivalent to  $\text{B}\Sigma_2^0$ .*

Notice that the proof in  $\text{B}\Sigma_2^0$  uses the finite van der Waerden's theorem.

### Corollary

*Over  $\text{RCA}_0$  (even  $\text{RCA}_0^*$ ), Brown's lemma implies van der Waerden's theorem.*

## A proof in $\text{I}\Sigma_3^0$

### Theorem

Over  $\text{RCA}_0$ ,  $\text{I}\Sigma_3^0$  implies Brown's Lemma.

### Proof.

Let  $f: \mathbb{N} \rightarrow r$  be a finite coloring. By bounded  $\Sigma_3^0$ -comprehension let:

$$C = \{c < r: \{x \in \mathbb{N}: f(x) \geq c\} \text{ is piecewise syndetic}\}.$$

Let  $c$  be the maximum element of  $C$ . If  $c + 1 = r$ , we are done. Suppose  $c + 1 < r$ .

Let  $d \in \mathbb{N}$  be such that  $\{x \in \mathbb{N}: f(x) \geq c\}$  is  $d$ -ps. By the choice of  $c$ ,  $\{x \in \mathbb{N}: f(x) \geq c + 1\}$  is not ps, and in particular is not  $d$ -ps. Then there must be an  $e \in \mathbb{N}$  such that every finite set of size  $e$  and gaps bounded by  $d$  contains an  $x$  such that  $f(x) \leq c$ . Show that  $\{x \in \mathbb{N}: f(x) = c\}$  is  $d \cdot e$ -ps.  $\square$

## From Brown's Lemma to $\text{I}\Sigma_2^0$

The number  $d$  does not depend on  $r$ .

### Lemma ( $\text{RCA}_0$ , Diagonalization Lemma)

*There exists a function  $g: \mathbb{N} \times \mathbb{N} \rightarrow 2$  such that for all  $d$  the finite 2-coloring  $f(x) = g(d, x)$  has no  $f$ -homogeneous  $d$ -piecewise syndetic sets.*

**Proof.**

For  $d = 3$  let 000111000111000...  $\square$

### Theorem

Over  $\text{RCA}_0$  (even  $\text{RCA}_0^*$ ), Brown's Lemma implies  $\text{I}\Sigma_2^0$ .

**Proof.**

We aim to prove  $\text{S}\Pi_1^0$ , strong collection for  $\Pi_1^0$ -formulas, that is:

$$(\forall a)(\exists b)(\forall n < a)(\exists m \theta(n, m) \rightarrow (\exists m < b)\theta(n, m)),$$

where  $\theta$  is  $\Pi_1^0$ .

Define  $h: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  such that  $h(n) = \lim_{s \rightarrow \infty} h(n, s) =$  the least  $m$  such that  $\theta(n, m)$ , provided that  $\exists m \theta(n, m)$ .

Suppose for a contradiction that there is an  $a$  with no  $b$ . Then for every  $b$  there is an  $n < a$  such that  $\exists m \theta(n, m)$  and  $h(n) \geq b$ .

Define a coloring  $f: \mathbb{N} \rightarrow 2^a$  as follows. Let  $g: \mathbb{N} \times \mathbb{N} \rightarrow 2$  be as in the Diagonalization Lemma and let  $f(x) = \langle g(h(n, x), x) : n < a \rangle$ .

By Brown's Lemma, there exists  $c \in 2^a$  such that  $\{x \in \mathbb{N} : f(x) = c\}$  is ps. Let  $e$  be a witness. For such an  $e$ , there exists  $n < a$  such that  $h(n)$  exists and  $d = h(n) \geq e$ .

Show that  $\{x \in \mathbb{N} : g(d, x) = c(n)\}$  is  $d$ -ps, for the desired contradiction. Let  $s$  be such that  $h(n, x) = h(n) = d$  for all  $x > s$ . Fix a size  $k$ . Let  $G$  be a finite set of size  $s + 1 + k$  and gaps bounded by  $e$  such that  $f(x) = c$  for all  $x \in G$ .

Then  $G$  contains at least  $k$  elements  $> s$ . Let  $F$  consist of the last  $k$  elements of  $G$ . Clearly  $F$  has size  $k$  and gaps bounded by  $d$ . Moreover, for all  $x \in F$  we have that

$$g(d, x) = g(h(n), x) = g(h(n, x), x) = f(x)(n) = c(n). \quad \square$$

# A proof in $\text{I}\Sigma_2^0$ .

## Definition ( $\text{RCA}_0$ )

Let  $H \subseteq \mathbb{N}$  be finite. Define the *gap size* of  $H$ , denoted  $gs(H)$ , as the largest difference between two consecutive elements of  $H$ . In other words, the gap size of  $H$  is the least  $d$  such that  $H$  has gaps bounded by  $d$ .

## Theorem (Brown's Lemma, finite)

Let  $k: \mathbb{N} \rightarrow \mathbb{N}$ . Then for all  $r$  there exists  $N$  such that every  $C: N \rightarrow r$  has a  $C$ -homogeneous set  $H$  of size  $\geq k(gs(H))$ .

## Theorem

The finite version of Brown's Lemma is provable in  $\text{RCA}_0$ .



## Lemma ( $\text{RCA}_0 + \text{I}\Sigma_2^0$ )

Let  $k: \mathbb{N} \rightarrow \mathbb{N}$  be a  $\Sigma_2^0$ -definable total function. Then for all  $r$  there exists  $N$  such that every  $P: \mathbb{N} \rightarrow r$  has a  $P$ -homogeneous set  $H$  of size  $\geq k(\text{gs}(H))$ .

## Theorem

Over  $\text{RCA}_0$ ,  $\text{I}\Sigma_2^0$  implies Brown's Lemma.

## Proof.

Let  $f: \mathbb{N} \rightarrow r$  be a finite coloring and suppose for a contradiction that  $\{x \in \mathbb{N}: f(x) = c\}$  is not piecewise syndetic for every  $c < r$ . Then for all  $d$  and  $c < r$  there exists a  $k$  such that no set of size  $k$  and gaps bounded by  $d$  is  $f$ -homogeneous for color  $c$ . By  $\text{B}\Sigma_2^0$ , for all  $d$  there exists  $k$  such that no set of size  $k$  and gaps bounded by  $d$  is  $f$ -homogeneous.

There exists a  $\Sigma_2^0$ -definable (in fact  $\Delta_2^0$ -definable) total function  $k: \mathbb{N} \rightarrow \mathbb{N}$  such that for all  $d$  no set of size  $k(d)$  and gaps bounded by  $d$  is  $f$ -homogeneous.

By the modified finite Brown's Lemma, let  $N$  be large enough be such that every  $C: N \rightarrow r$  has a  $C$ -homogeneous set  $H$  of size  $\geq k(gs(H))$ . In particular there exists an  $f$ -homogeneous set  $H$  of size  $\geq k(gs(H))$ , for the desired contradiction.  $\square$

## Questions

Let  $B(r, k)$ , for  $r > 0$  and  $k: \mathbb{N} \rightarrow \mathbb{N}$ , be the least natural number  $N$  such that every  $r$ -coloring of  $N$  has a homogeneous set  $F$  of size at least  $k(\text{gs}(F))$ . The proof of the finite Brown's lemma gives superexponential upper bounds  $N(r, k)$  for  $B(r, k)$ . For instance, if  $k(d) = 2^d$ , then  $N(r, k) \geq 2_r$ , where

$$2_r = 2^{2^{\dots^2}}, \text{ a tower of } r \text{ twos.}$$

### Question

*Is finite Brown's lemma provable in  $\text{RCA}_0^*$ ?*

We expect the answer to be positive.

As the finite van der Waerden's theorem is already provable in  $\text{RCA}_0$  and presumably in  $\text{RCA}_0^*$ , the question whether Brown's lemma implies van der Waerden's theorem and vice versa can be settled only over a weak system of arithmetic.

### Question

*What is the relationship between (the finite versions of) Brown's lemma and van der Waerden's theorem over a suitable bounded fragment of second-order arithmetic?*

## Question (Mummert)

*What is the Weihrauch degree of Brown's lemma?*

$BL \subseteq \{f: \mathbb{N} \rightarrow \mathbb{N}\} \rightrightarrows \mathbb{N}$ , where

$BL(f) = \{c: f^{-1}(c) \text{ is piecewise syndetic}\}$ .

$BL_r \subseteq \{f: \mathbb{N} \rightarrow r\} \rightrightarrows \mathbb{N}$ , where

$BL_r(f) = \{c < r: f^{-1}(c) \text{ is piecewise syndetic}\}$ .

$BL_r^* \subseteq \{f: \mathbb{N} \rightarrow r\} \rightrightarrows \mathbb{N}^2$ , where

$BL_r^*(f) = \{(c, d): f^{-1}(c) \text{ is } d\text{-piecewise syndetic}\}$ .

- $RT_r^1 \leq_{sW} BL_r$ .
- $C_{\mathbb{N}} \leq_{sW} BL_2$  (Mummert).

Tks