

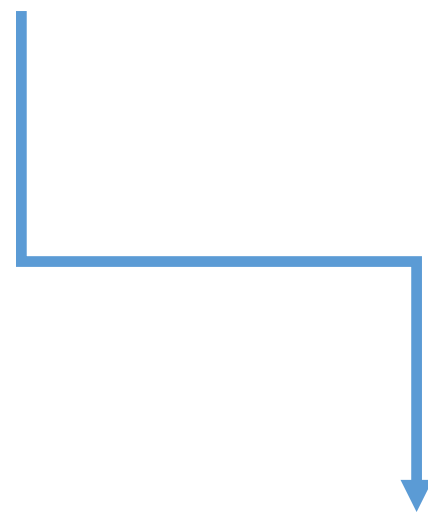
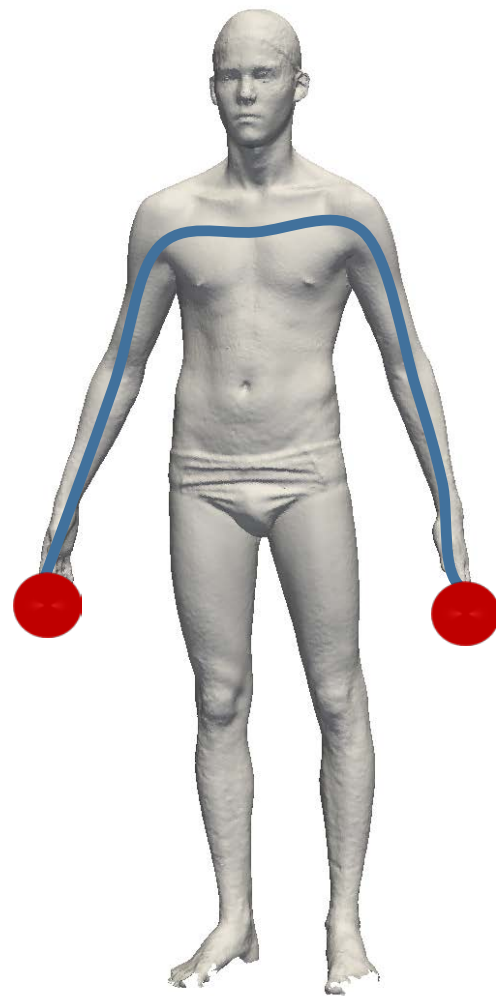
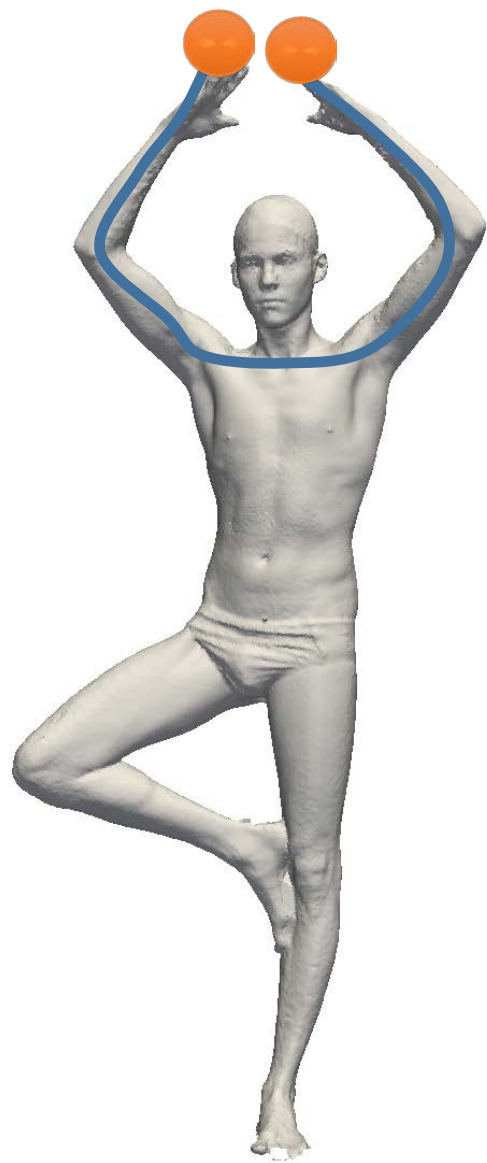
# On the Applicability of Convex Relaxations for Matching Symmetric Shapes

Nadav Dym

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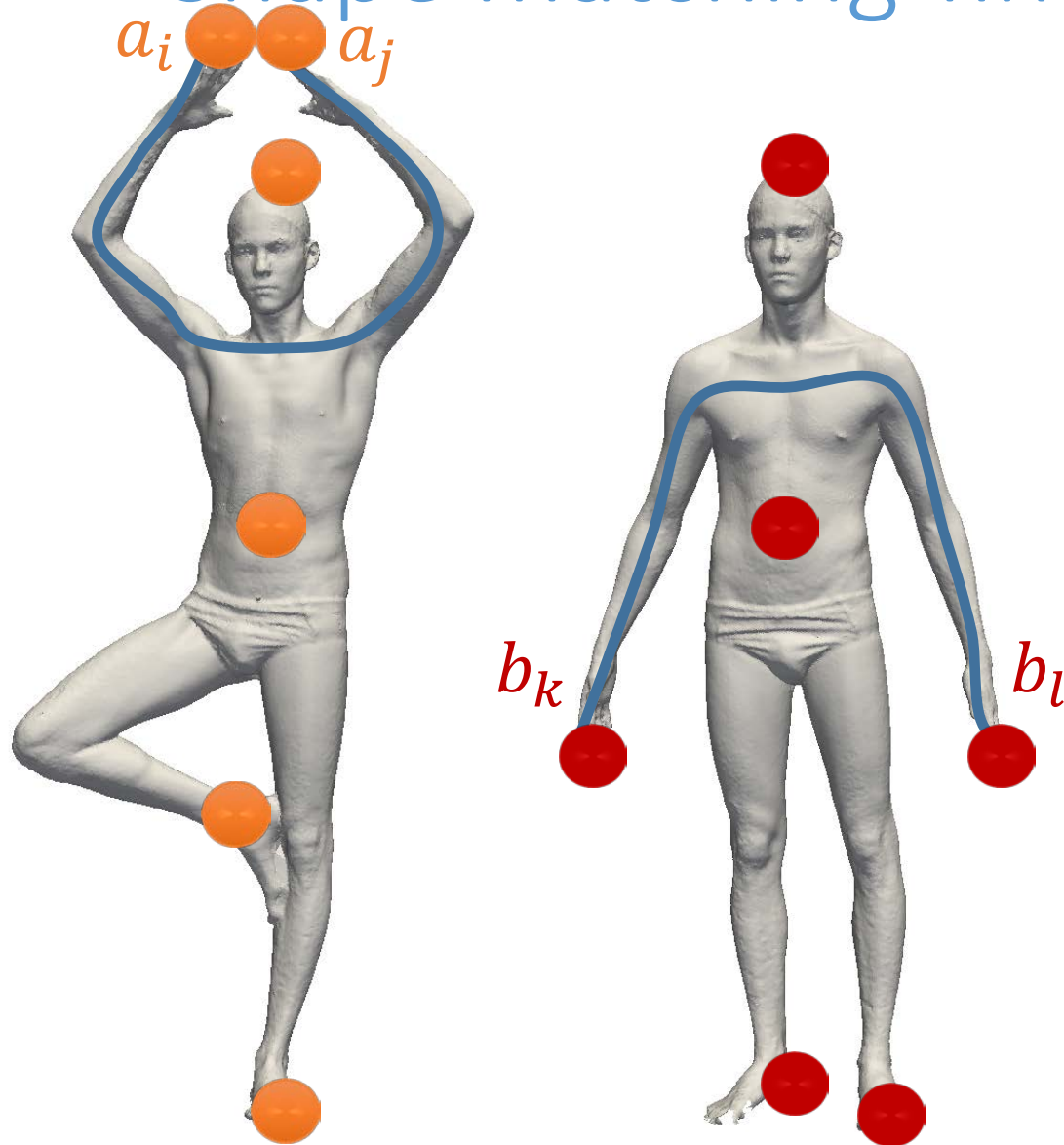
In collaboration with *Haggai Maron* and *Yaron Lipman*

# Shape matching



Finding isometries

# Shape matching-finding isometries



$$A_{ij} = d(a_i, a_j)$$

$$B_{kl} = d(b_k, b_l)$$

Goal: find mapping (permutation)

$\sigma: \{a_1, \dots, a_n\} \rightarrow \{b_1, \dots, b_n\}$   
such that

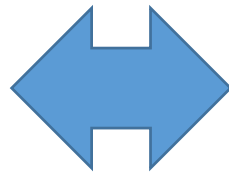
$$A_{ij} = B_{\sigma(i), \sigma(j)}$$

# Graph isomorphism

Input:  $A = A^T, B = B^T$

Goal: find permutation (if exists)  $\sigma$  such that

$$A_{ij} = B_{\sigma(i),\sigma(j)}$$



Goal: find permutation matrix  $P \in \Pi_n$  such that

$$A = P^T B P$$

# Graph matching/quadratic assignment

Input:  $A = A^T, B = B^T$

Output:  $P \in \Pi_n$  such that  $A \approx P^T B P$ :

$$P^* = \operatorname{argmin}_{P \in \Pi_n} \|A - P^T B P\|_F$$

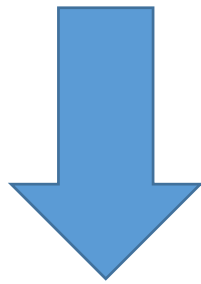
$$P^* = \operatorname{argmin}_{P \in \Pi_n} \|PA - BP\|_F$$

Graph matching

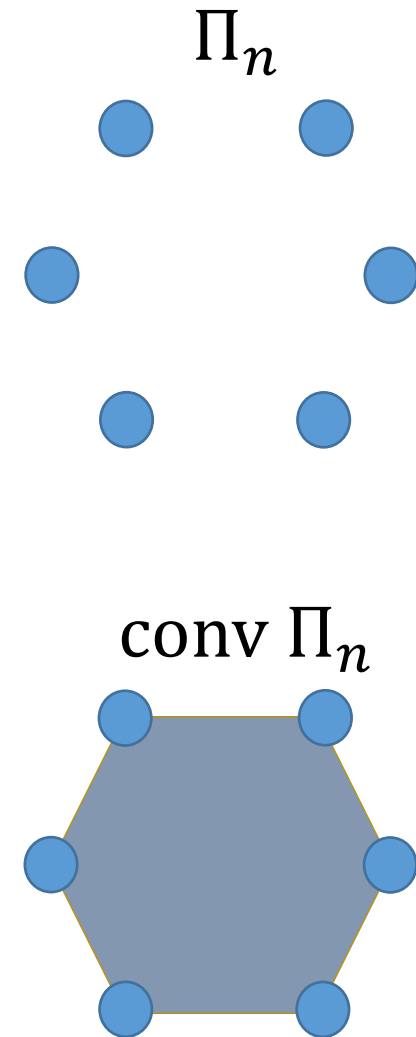
# DS relaxation

$$P^* = \operatorname{argmin}_{P \in \Pi_n} \|PA - BP\|_F$$

NP Hard!



$$S^* = \operatorname{argmin}_{S \in \operatorname{conv} \Pi_n} \|SA - BS\|_F$$



# DS relaxation

$\text{conv } \Pi_n$  is the Birkhoff polytope:

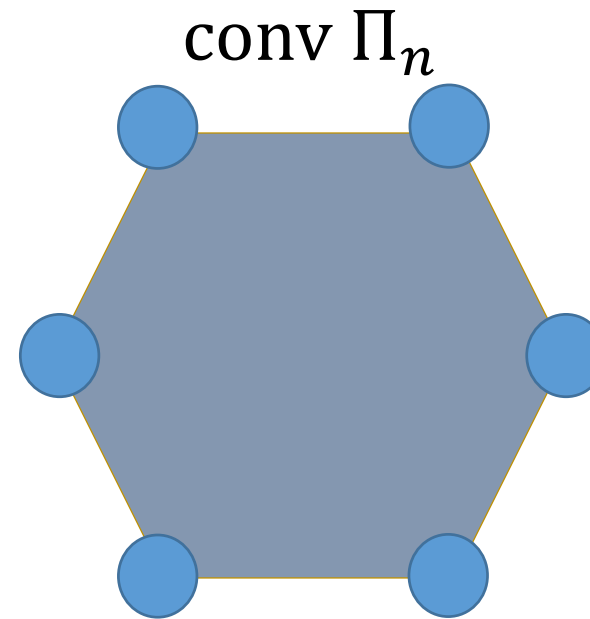
$$DS = \{S \mid S \geq 0, S\mathbf{1} = \mathbf{1}, S^T\mathbf{1} = \mathbf{1}\}$$

*DS relaxation:*

$$S^* = \underset{S \in DS}{\text{argmin}} \|SA - BS\|_F$$

↓

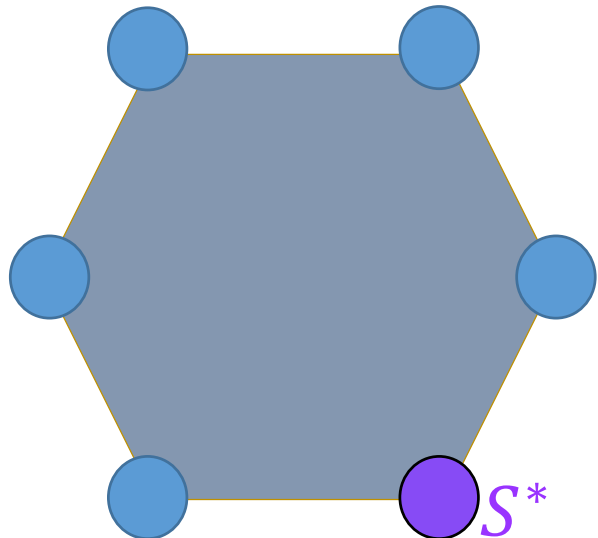
$$GM_{DS}(A, B)$$



# DS relaxation-does it work?

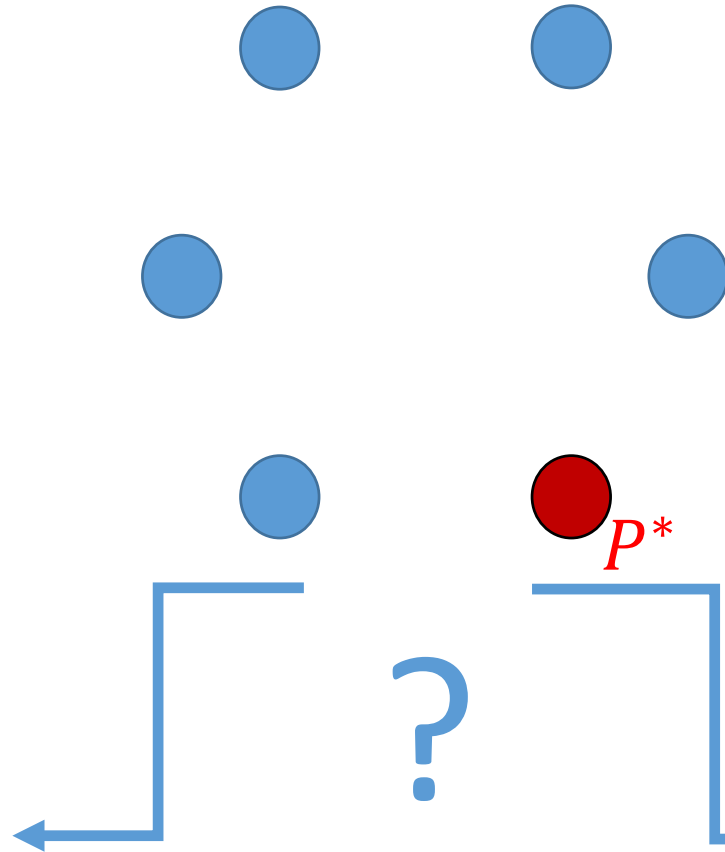
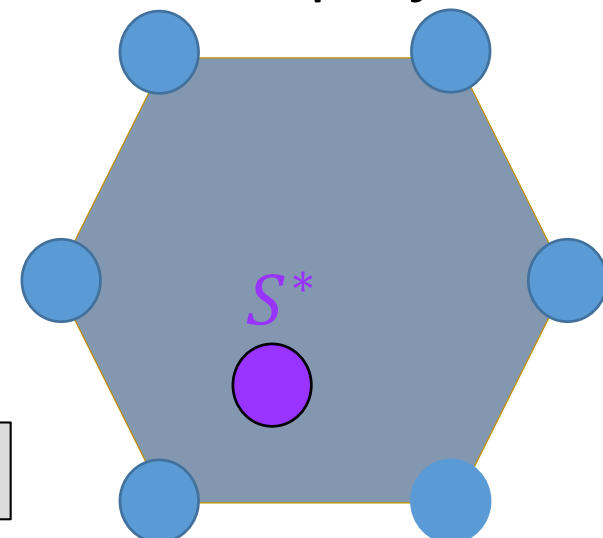
$$P^* = \operatorname{argmin}_{P \in \Pi_n} \|PA - BP\|_F$$

Part I: exactness



$$S^* = \operatorname{argmin}_{S \in DS} \|SA - BS\|_F$$

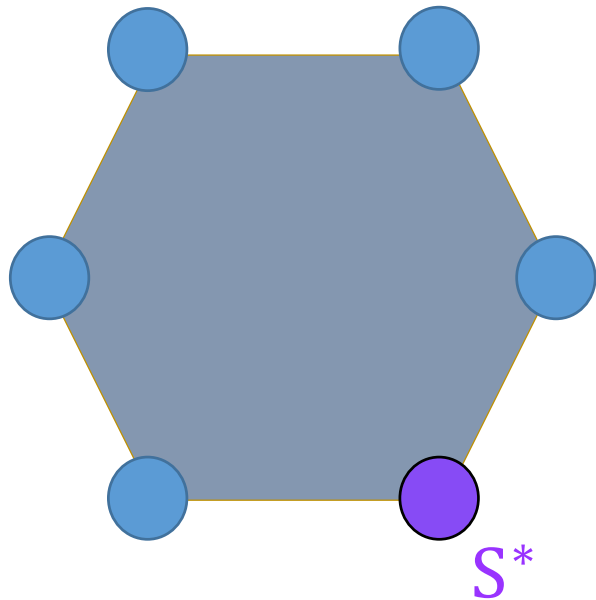
Part II: projection



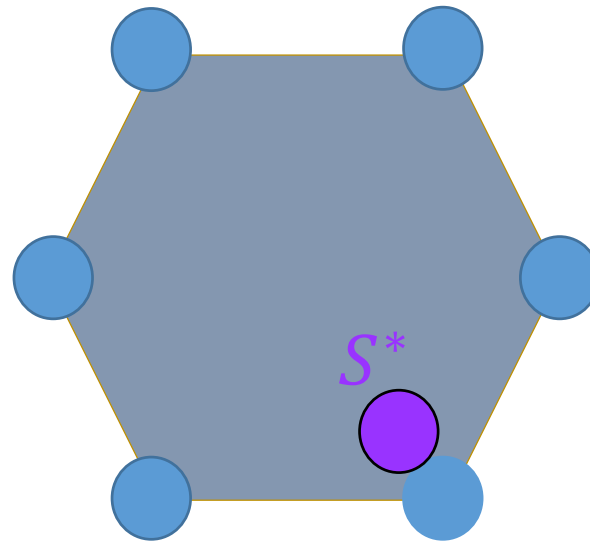


# Exactness affects projection efficiency

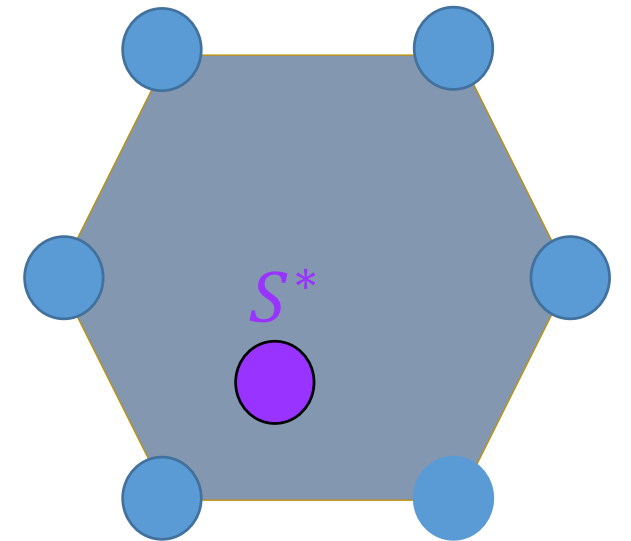
Exact case



Exact+small noise



Very noisy



$$S^* = \operatorname{argmin}_{S \in DS} \|SA - BS\|_F$$

# Part I: Exactness

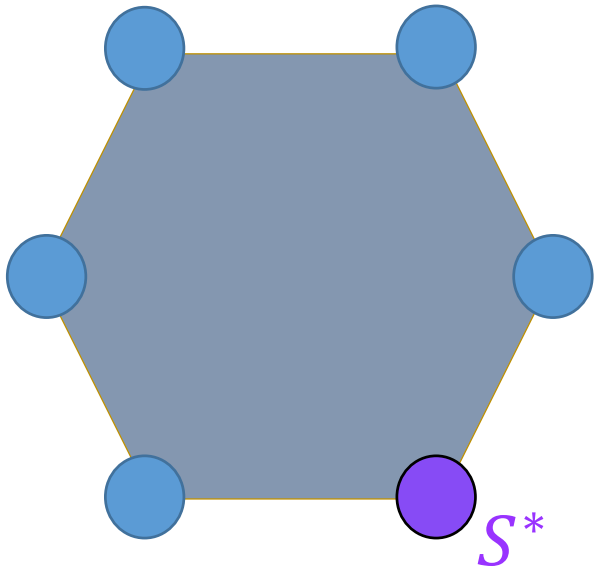
[Aflalo et al 2015], [Fiori and Sapiro 2015]

Assume: (i)  $A \cong B$  (ii) Unique isomorphism  $P^*$

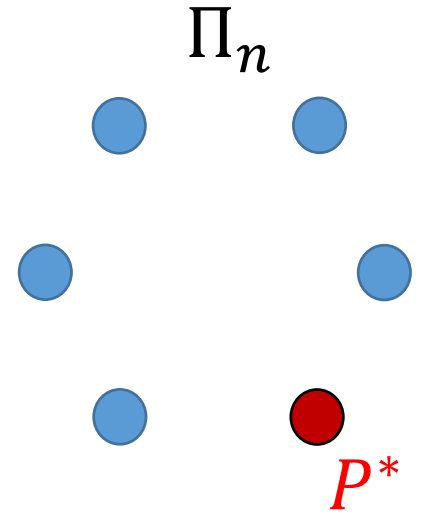
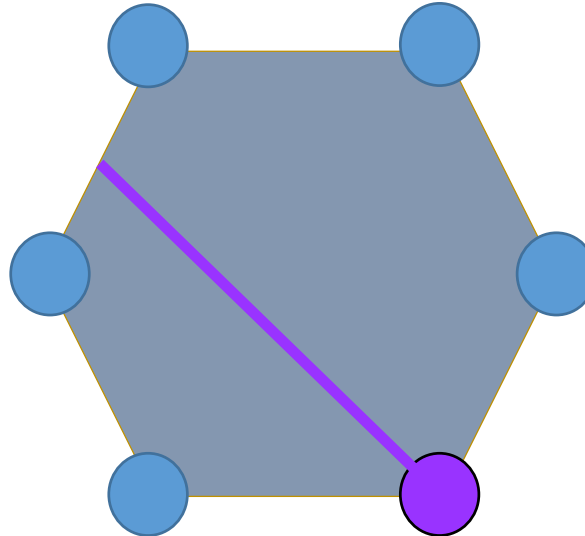


$$\min_{S \in DS} \|AS - SB\|_F = 0$$

Usually

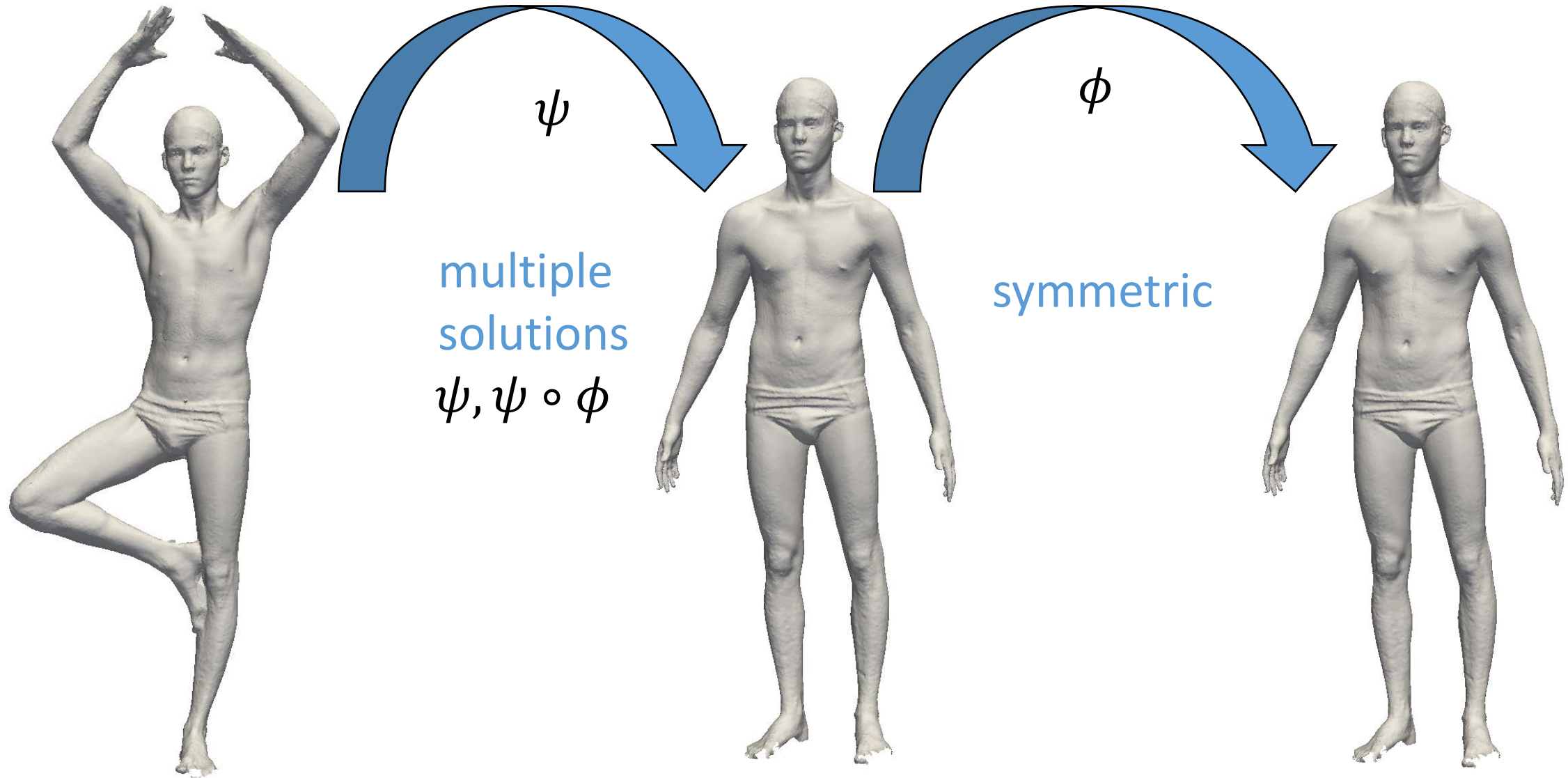


Usually not



$$S^* = \operatorname{argmin}_{S \in DS} \|SA - BS\|_F$$

# Problem: Unique solution assumption



# Symmetries of natural shapes

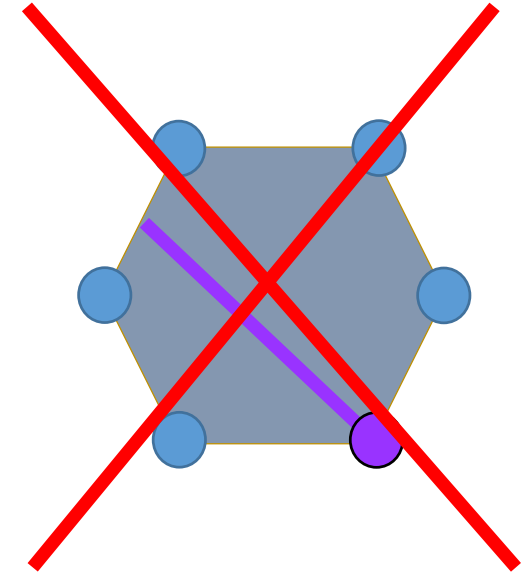
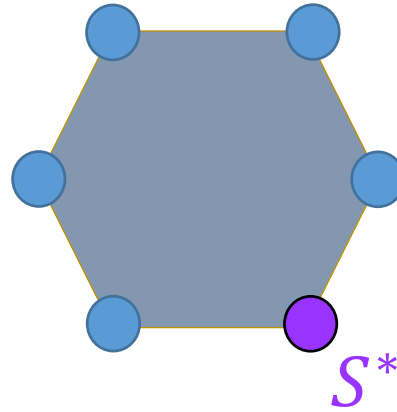
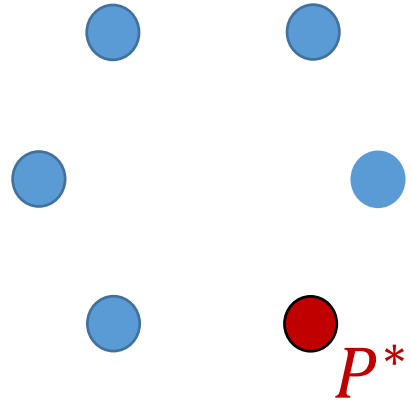


Bilateral symmetry:  
[SCAPE, FAUST, TOSCA]

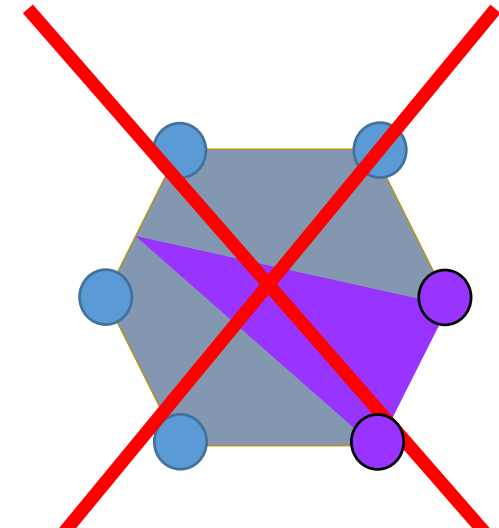
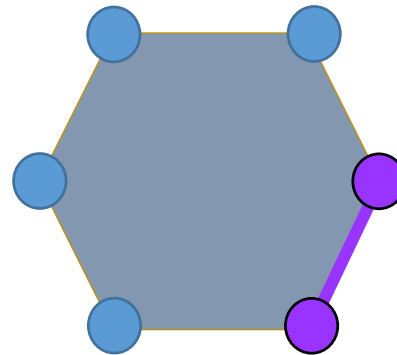
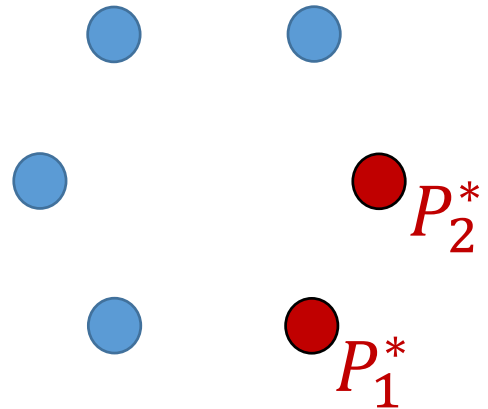
[SHREC]

# Exactness vs. convex exactness

Exactness (asymmetric)



Convex exactness (symmetric)

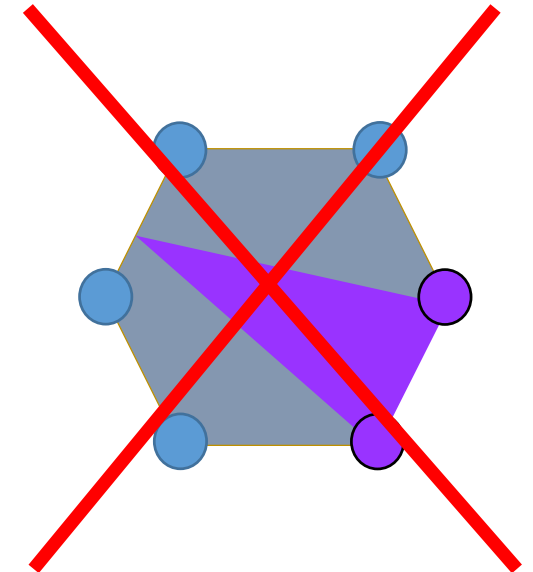
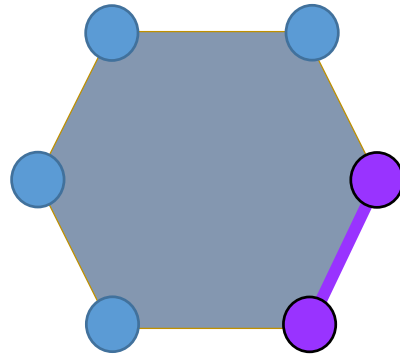
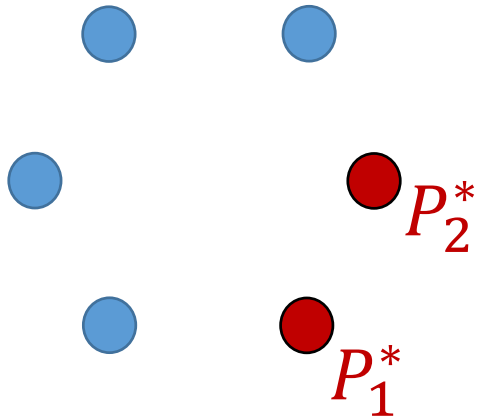


# Convex exactness-definition

$$Iso(A, B) = \{P \in \Pi_n \mid AP = PB\}, \quad Iso_{conv}(A, B) = \{S \in DS \mid AS = SB\}$$

$GM_{DS}(A, B)$  is convex exact if

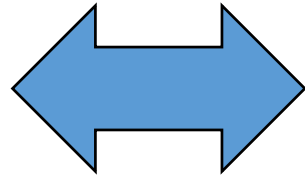
$$Iso_{conv}(A, B) = conv(Iso(A, B))$$



# A convenient reduction ( $B=A$ )

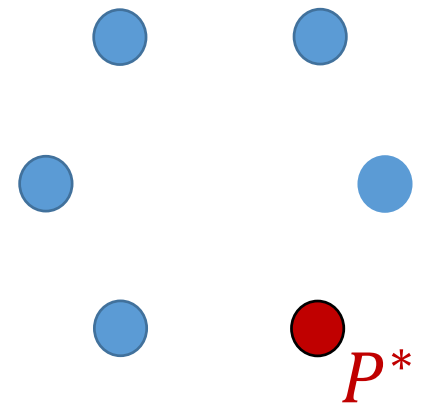
(easy) Lemma:

$GM_{DS}(A, A)$  is convex exact



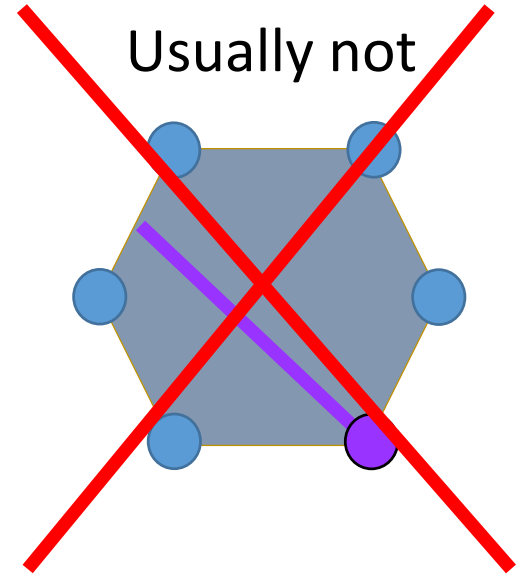
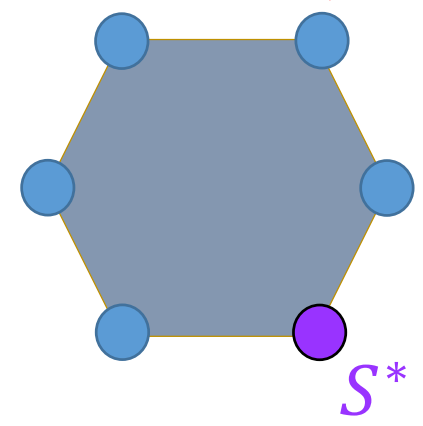
For any  $B$  s.t.  $B \cong A$ ,  
 $GM_{DS}(A, B)$  is convex exact

# Goal

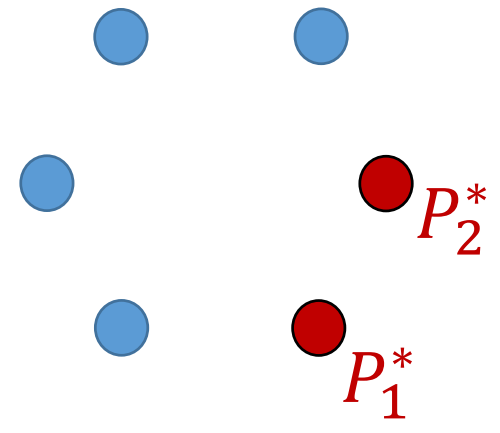


Almost surely

~~Usually~~

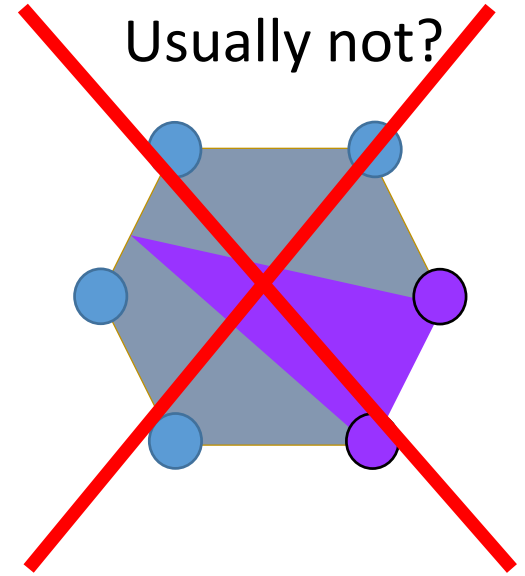
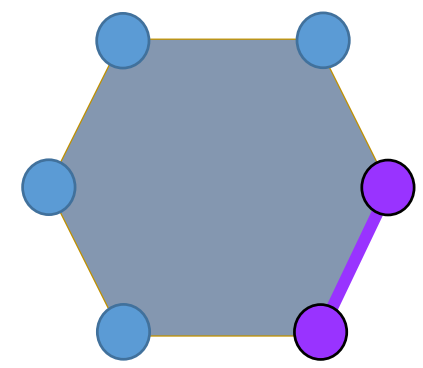


Usually not



Almost surely

~~Usually?~~



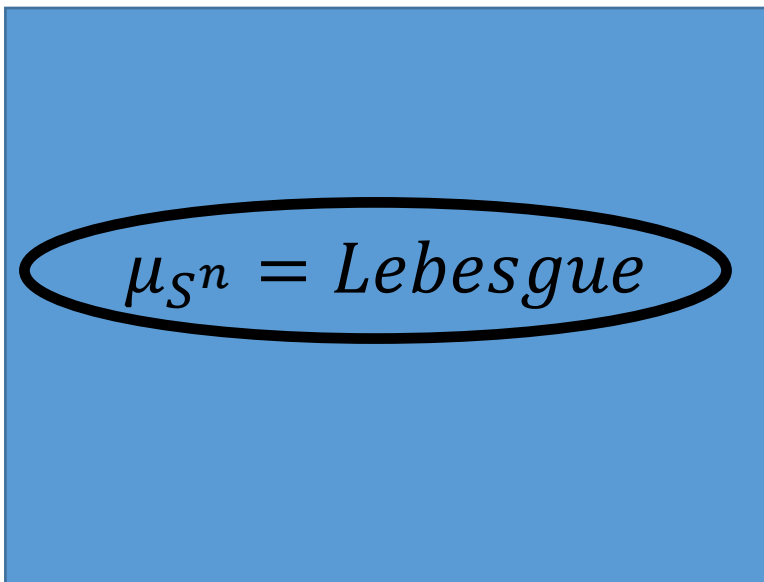
Usually not?



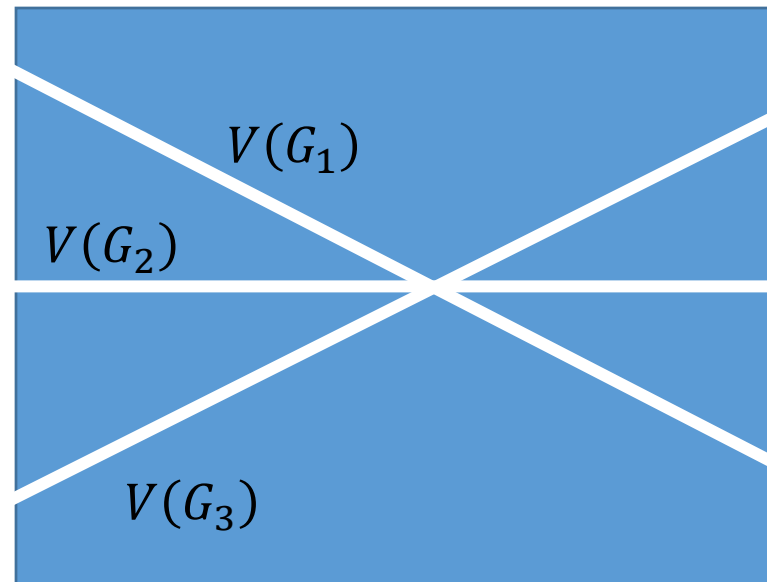
# Measure for the space of asymmetric graphs

Asymmetric graphs:  $\{A = A^T \mid A \text{ has no non-trivial automorphisms}\}$

$$S^n = \{A \mid A = A^T\}$$



Asymmetric graphs



$$G_i > \{I_n\}$$

$$V(G) = \{A \in S^n \mid PA = AP, \forall P \in G\}$$

# Measure for Graphs with prescribed symmetry group $G_0$

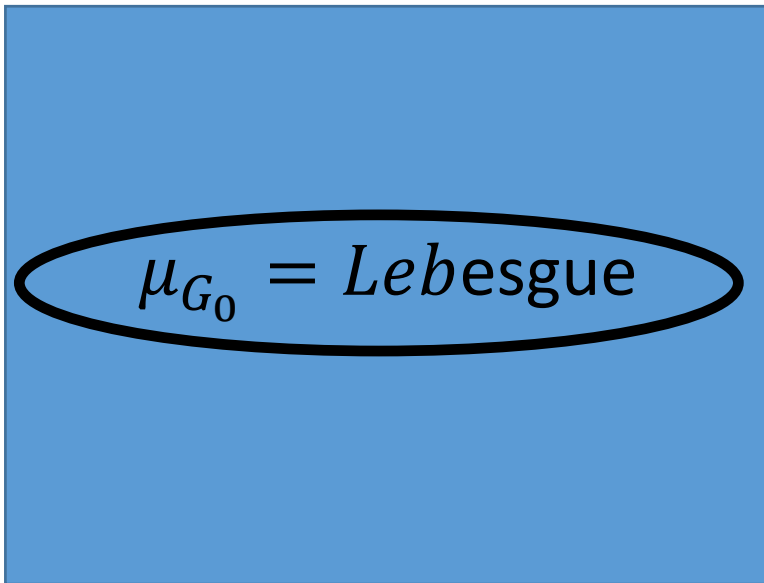
Graphs with sym group  $G_0$ : Graphs whose automorphism group is  $G_0$

In

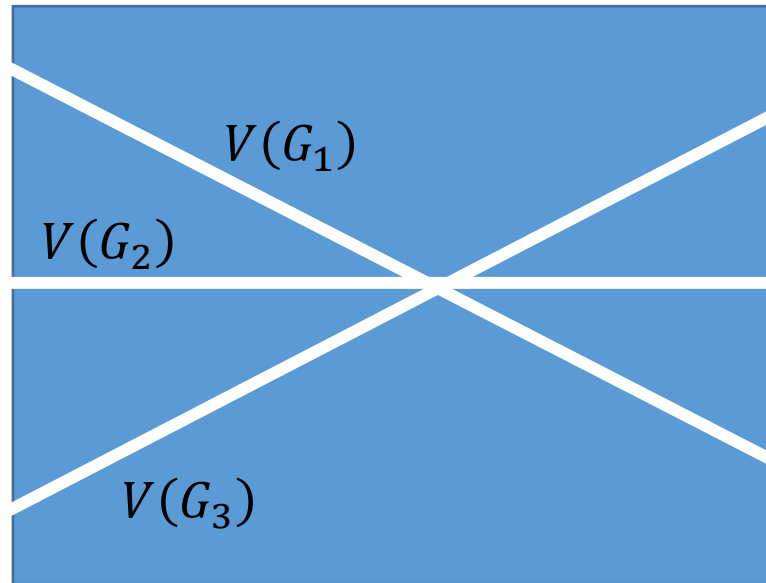
$$V(G_0) = \{A \in S^n \mid PA = AP, \forall P \in G_0\}$$

# Measure for Graphs with prescribed symmetry group $G_0$

$V(G_0)$



Graphs with sym group  $G_0$



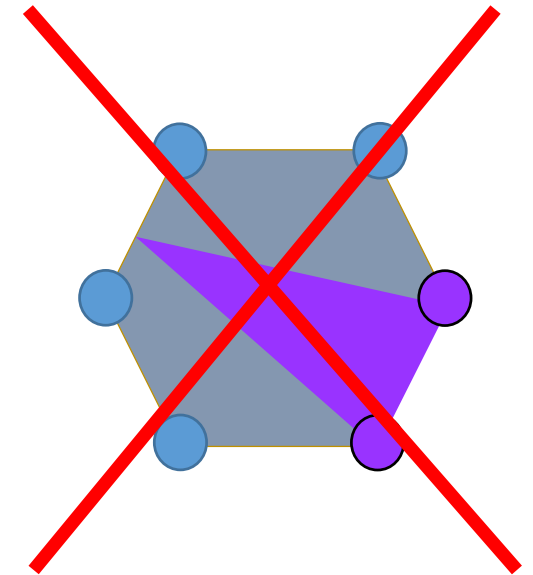
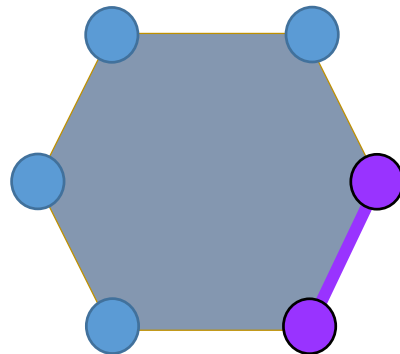
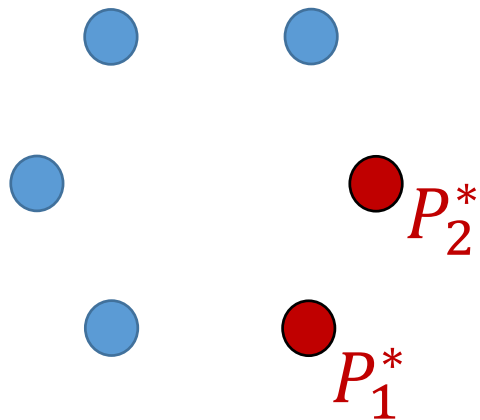
$G_i > G_0$

$$V(G) = \{A \in S^n \mid PA = AP, \forall P \in G\}$$

# Convex exactness for reflective groups

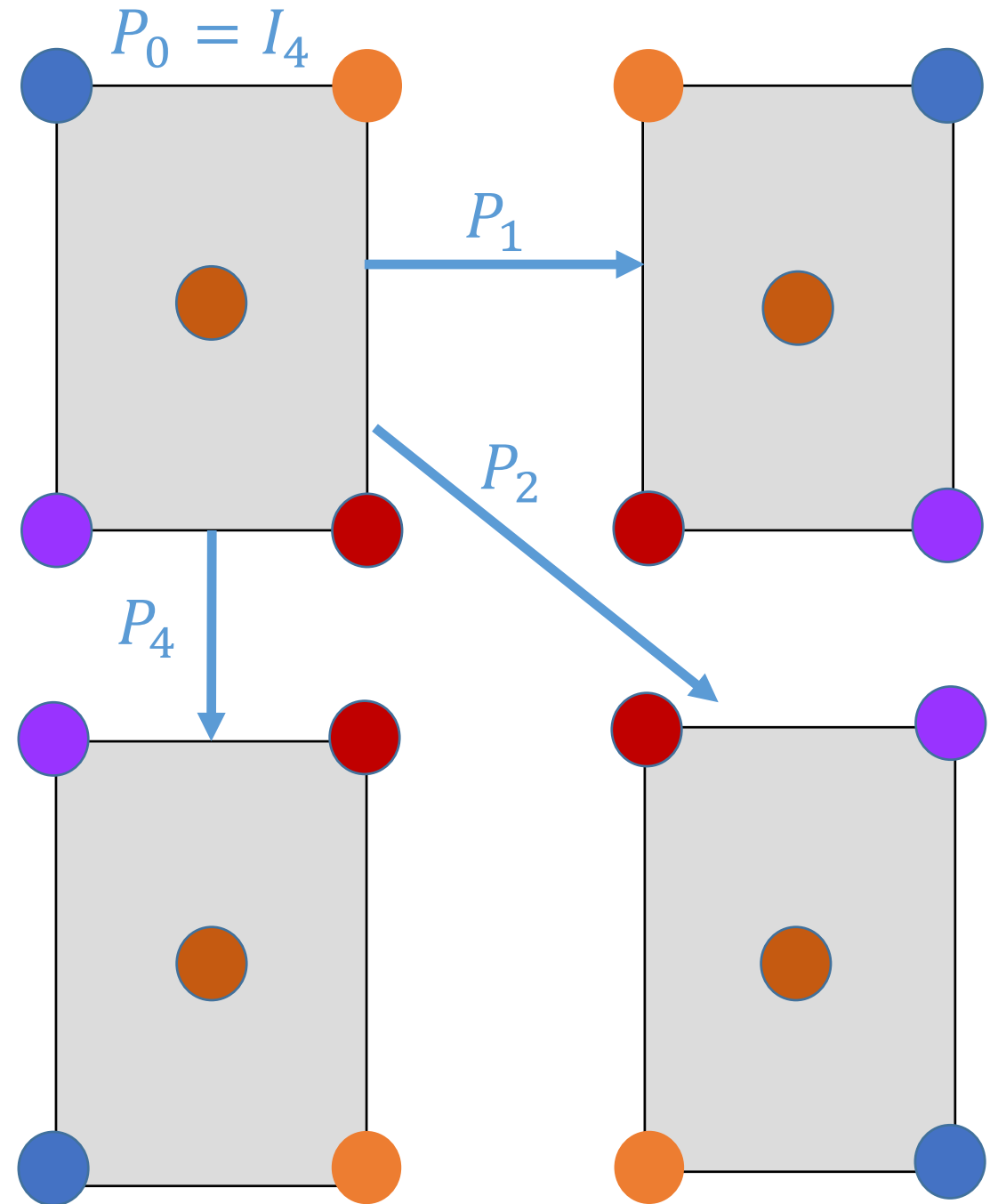
Theorem 1: If  $G \cong Z_2$ , then for  $\mu_G$  almost every  $A$ ,  $GM_{DS}(A, A)$  is convex exact.

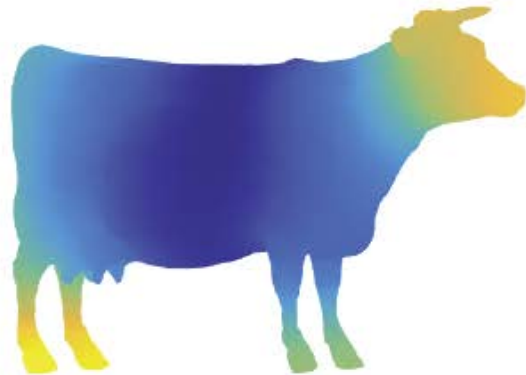
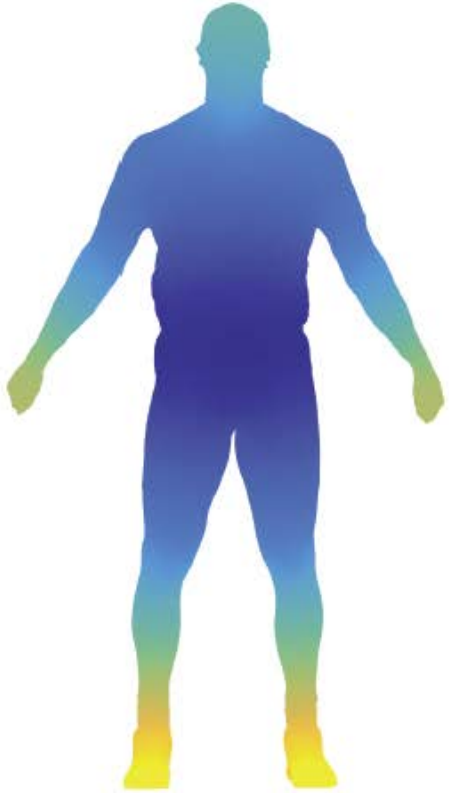
Also true for  $G = \{I_n\}$



In general Theorem 1 holds if

- $G$  is reflective ( $P^2 = I_n, \forall P \in G$ ).
- $G$ 's action on the vertices has a full orbit.





# General groups: 0-1 probability

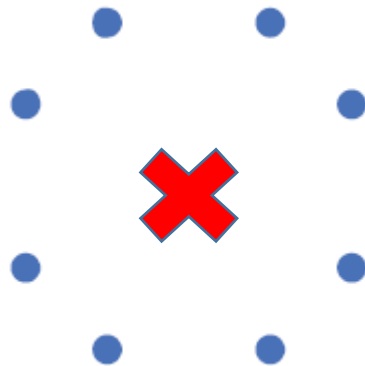
Theorem 2: For any  $G \leq \Pi_n$ , either

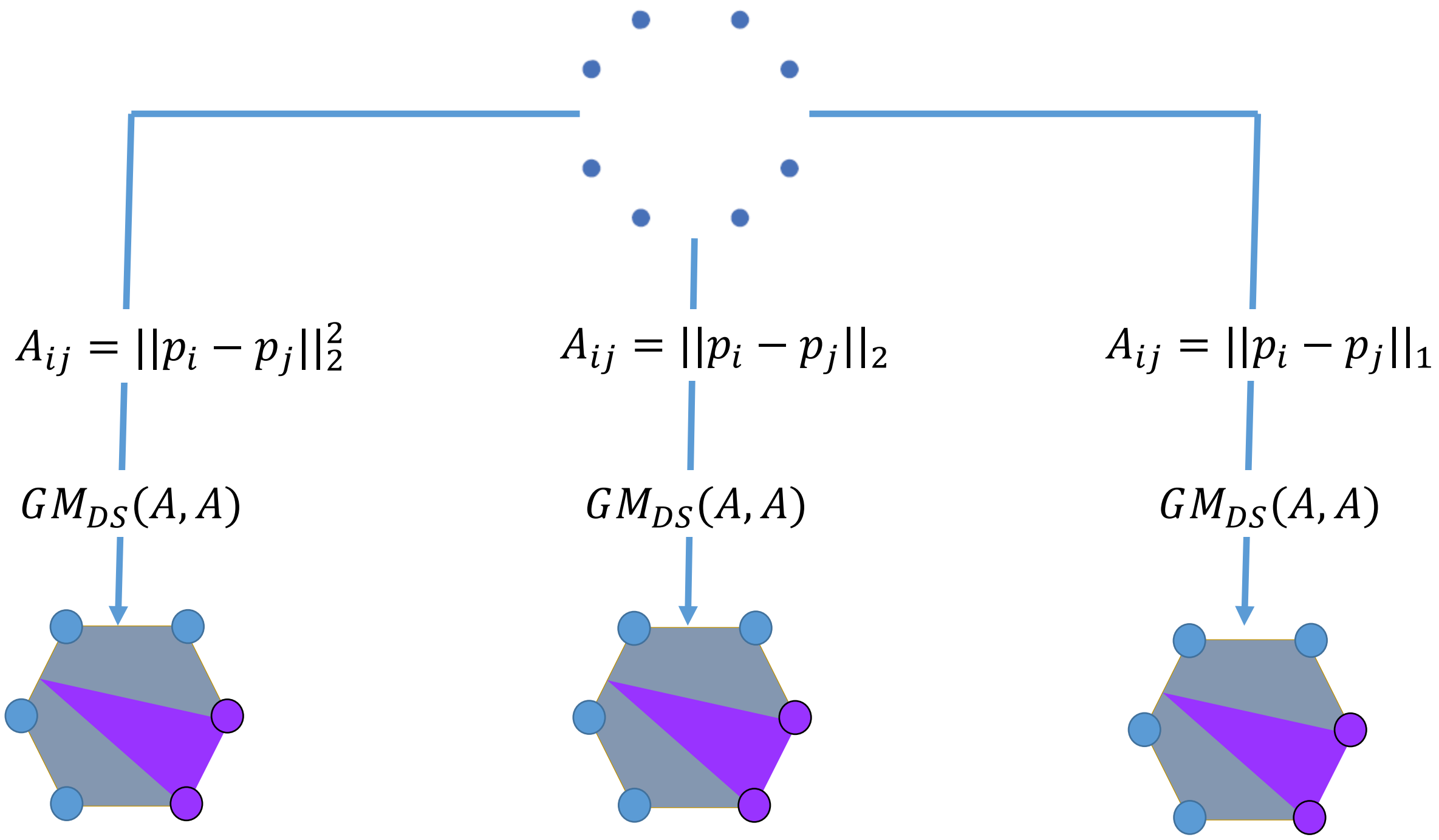
(i) For  $\mu_G$  a.e.  $A$ ,  $GM_{DS}(A, A)$  is convex exact.

Or

(ii) For **all**  $A$  with sym group  $G$ ,  $GM_{DS}(A, A)$  is **not** convex exact.

Proof is constructive...







# General groups: 0-1 probability

Theorem 2: For any  $G \leq \Pi_n$ , either

(i) For  $\mu_G$  a.e.  $A$ ,  $GM_{DS}(A, A)$  is convex exact.

Or

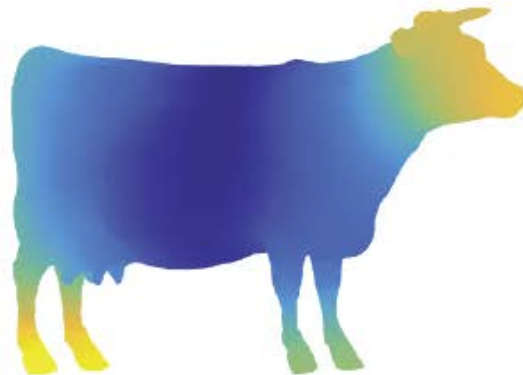
(ii) For **all**  $A$  with sym group  $G$ ,  $GM_{DS}(A, A)$  is **not** convex exact.



# Summary- Part I convex exactness



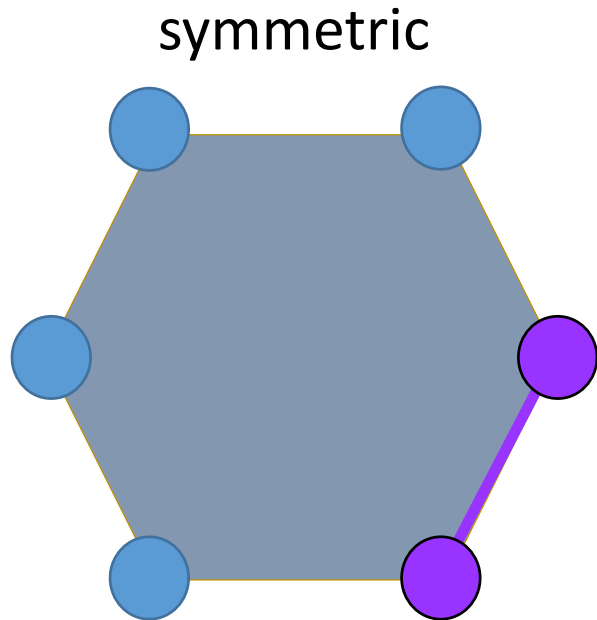
almost everywhere



almost everywhere



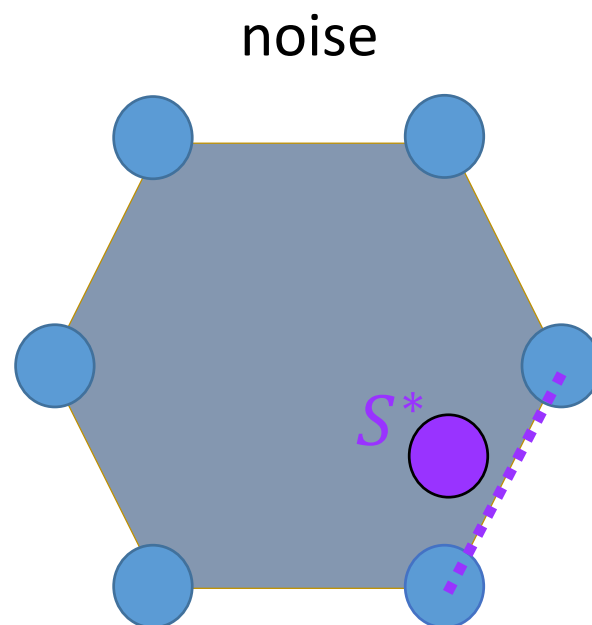
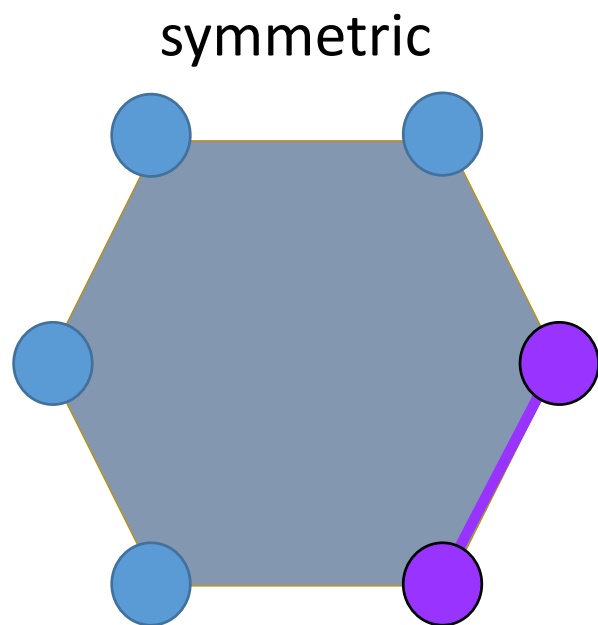
## Part II: Where's my permutation?



$$\text{set of minimizers} = \{S \in DS \mid AS = SB\}$$

↓  
Simplex  
algorithm

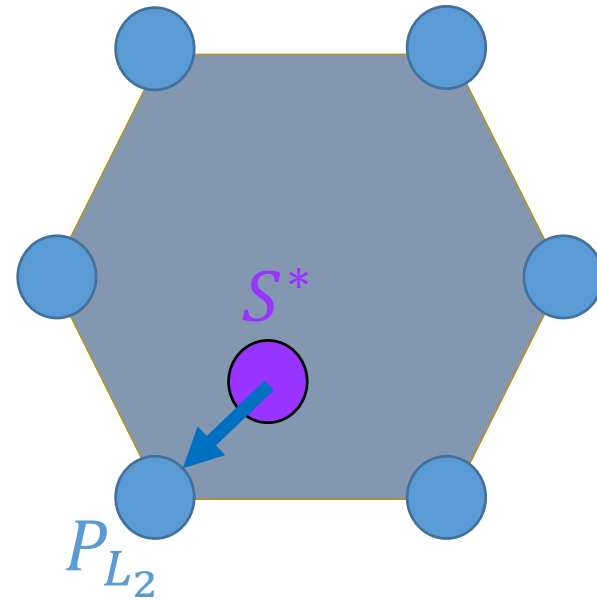
## Part II: Where's my permutation?



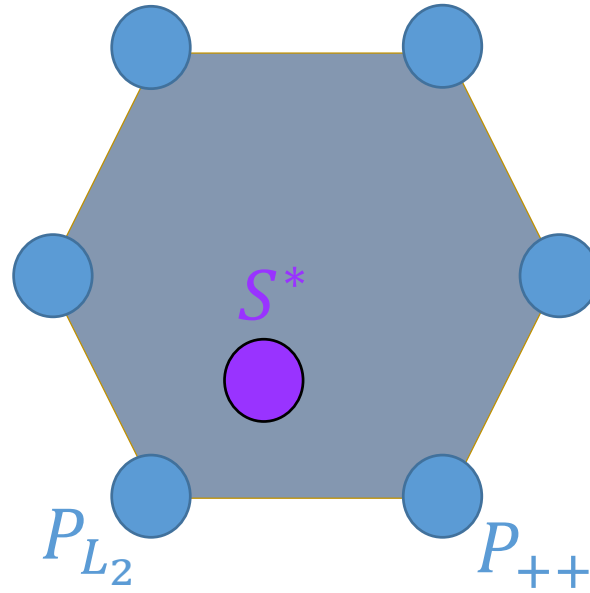
# DS relaxation- $L_2$ projection

$$S^* = \operatorname{argmin}_{S \in DS} \|SA - BS\|_F$$

$$P_{L_2} = \operatorname{argmin}_{P \in \Pi} \|S^* - P\|_F$$



# DS++: convex2concave projection



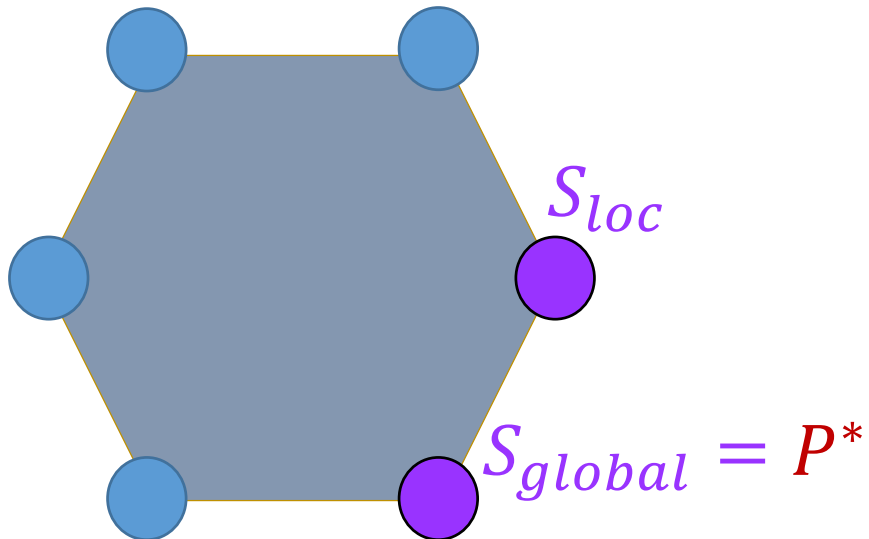
# convex2concave projection

[Zaslavskiy, Bach and Vert 2009]

Observation: Convex energy  $E_0$  is equivalent over  $\Pi_n$  to concave energy  $E_T$ .

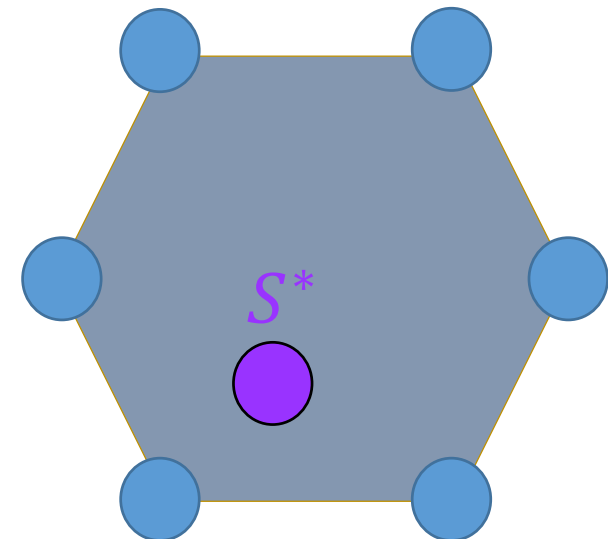
Concave energy:

- (i) Local/global minima are permutations!
- (ii) intractable



Convex energy:

- (i) minima may not be permutations
- (ii) tractable!



# convex2concave projection

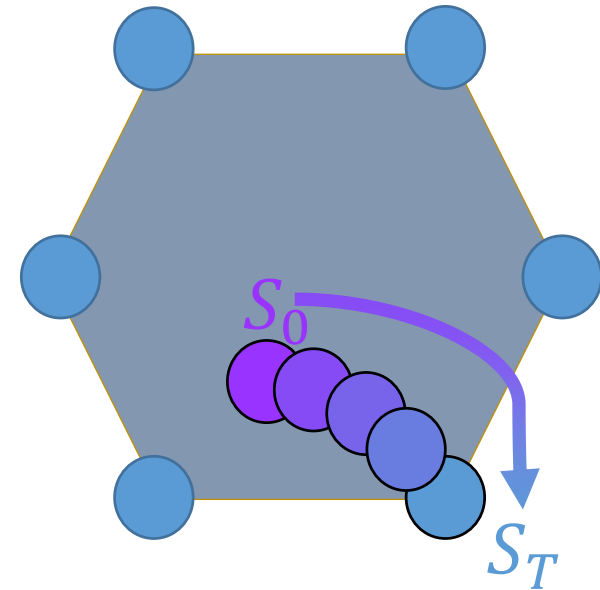
convex

concave

$E_{t_0}$   $E_{t_1}$   $E_{t_2}$  ...

$E_T$

$S_k = \text{"argmin"}_{S \in DS} E_{t_k}(S)$   
Warm start optimization from  $S_{k-1}$ .





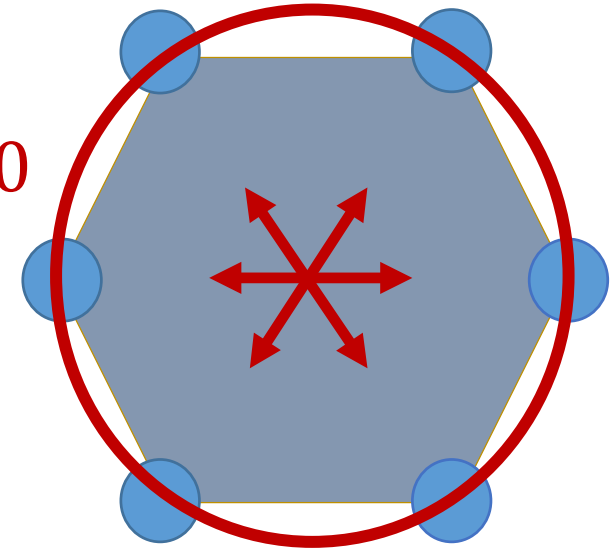
# DS++: convex2concave projection

$$E_t(S) = \|SA - BS\|_F^2 + t(n - \|S\|_F^2)$$

convex concave  
 $E_{t_0}$   $E_{t_1}$   $E_{t_2}$  ...  $E_{t_F}$

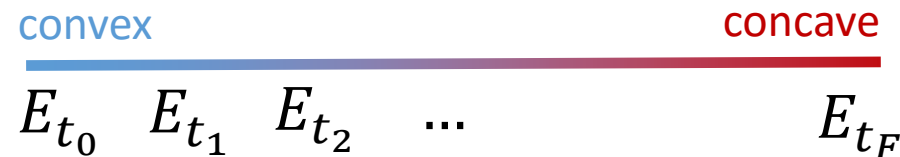
- $E_0 = E$
- $E_t = E$  over  $\Pi_n$
- $E_{t_F}$  strictly concave for  $t_F \gg 0$

$$n - \|S\|_F^2 = 0$$



# Choosing $[t_0, t_F]$

$$E_t(S) = \|SA - BS\|_F^2 + t(n - \|S\|_F^2)$$



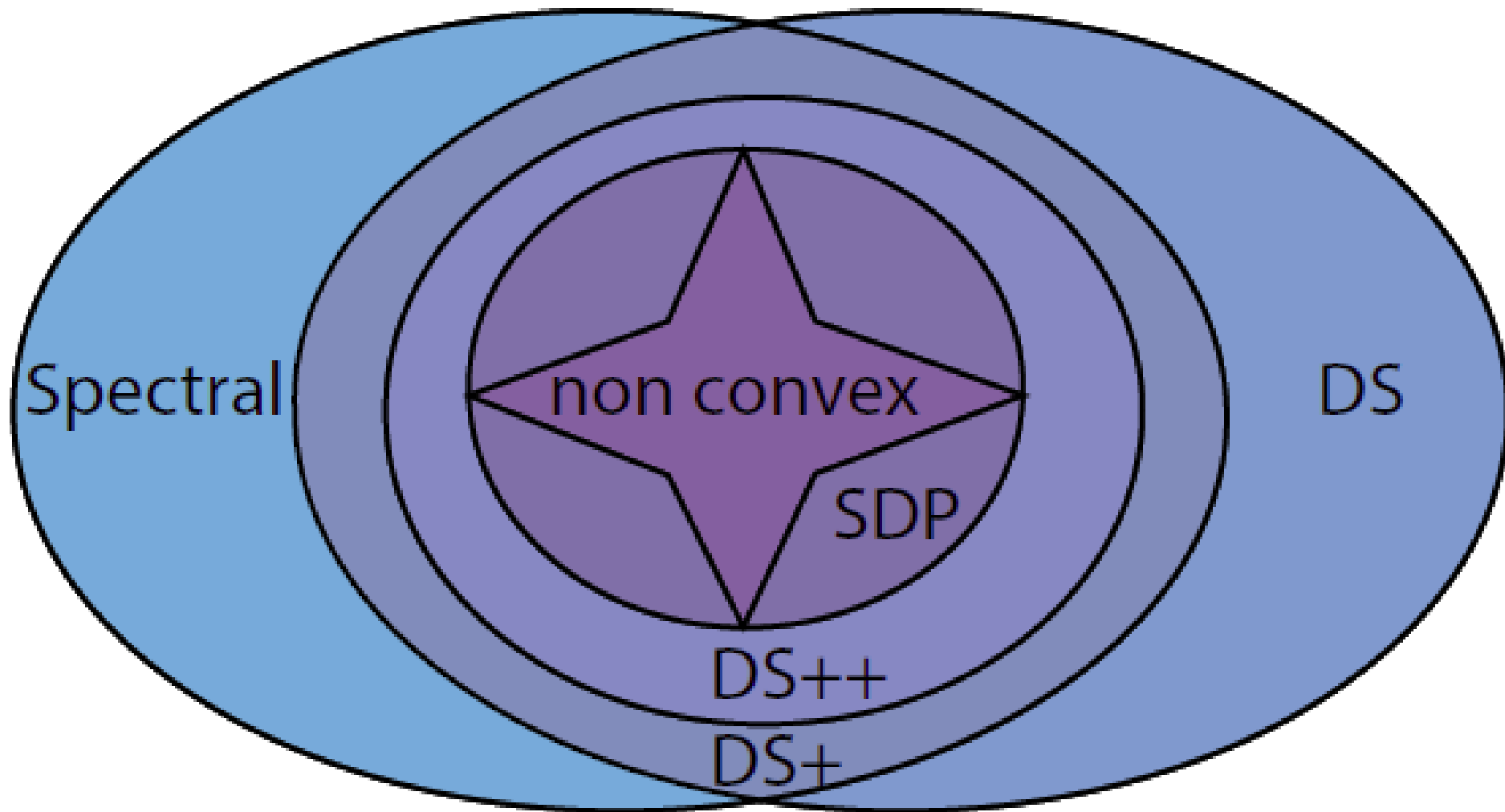
Best choice of  $t_F$ :

$$t_F = \lambda_{max}$$
$$t_F = \lambda_{max} \text{ over } V_{DS} = \{S \mid S\mathbf{1} = 0, S^T\mathbf{1} = 0\}$$

Best choice of  $t_0$ :

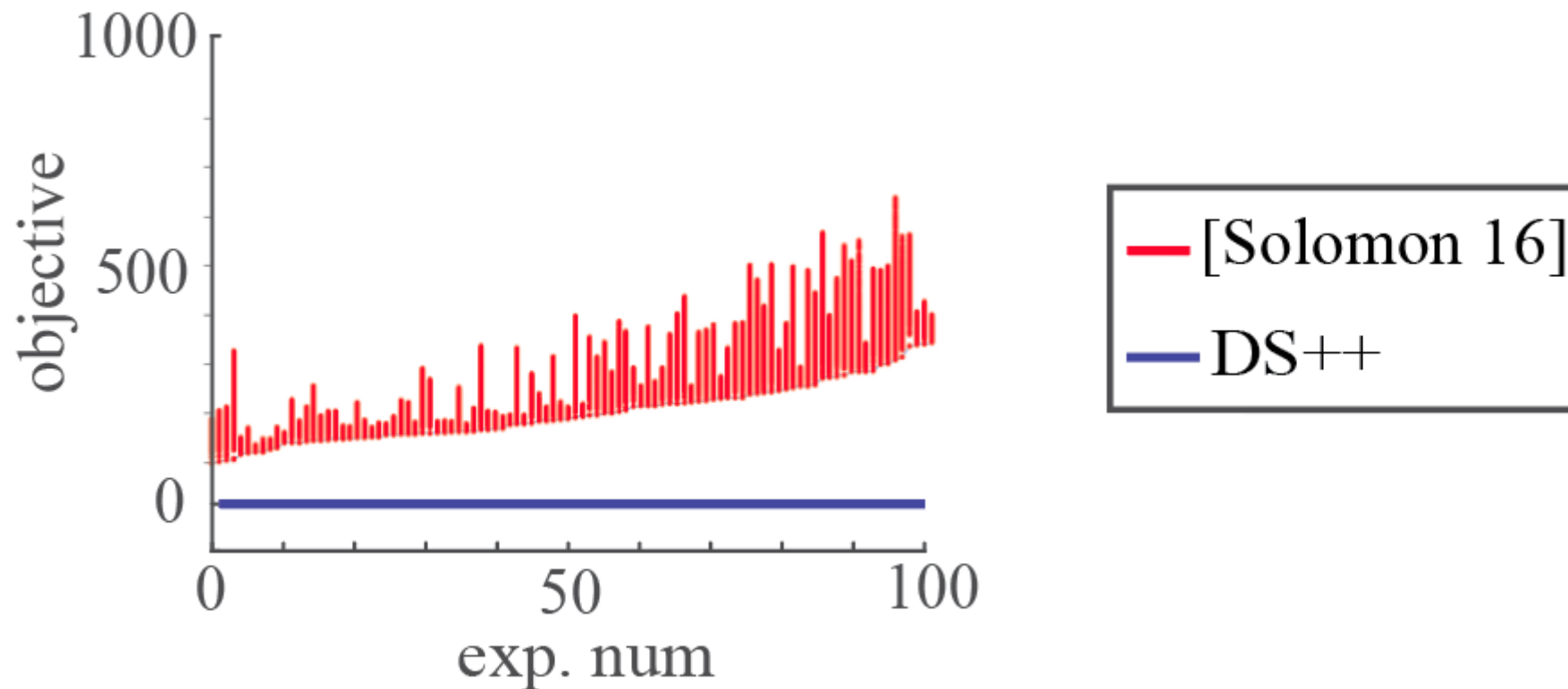
(DS)	$t_0 = 0$	[Aflalo et al. 15]
(DS+)	$t_0 = \lambda_{min}$	[Fogel et al. 13,15]
(DS++)	$t_0 = \lambda_{min} \text{ over } V_{DS}$	

# Relaxation comparison

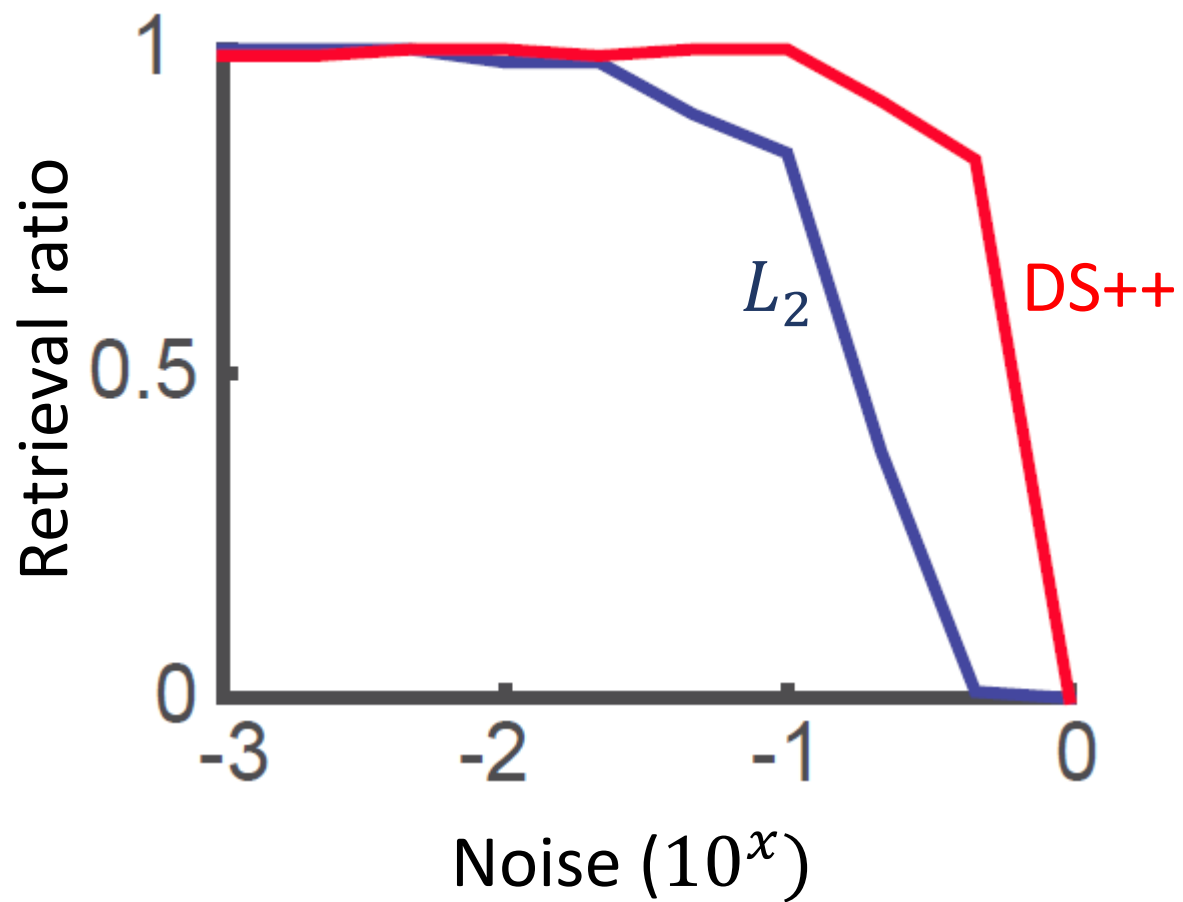


# DS++ vs local minimization

DS++ vs local minimization with 1000 different initializations:



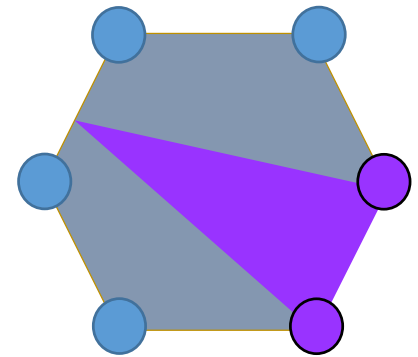
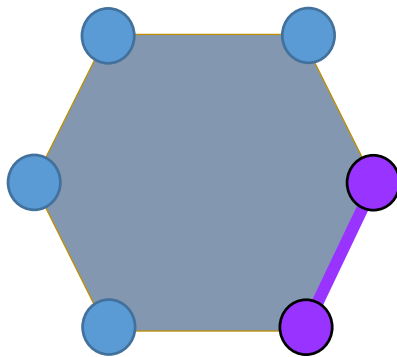
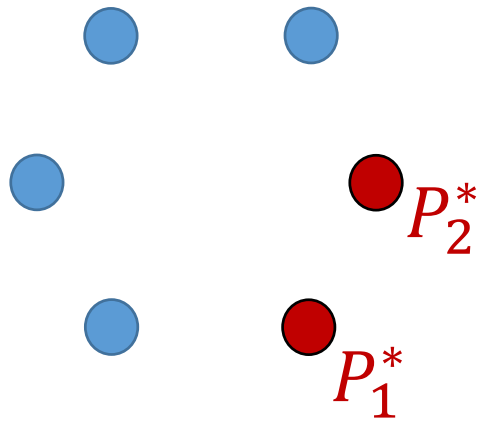
# Projection comparison



# Symmetric, no noise

$$E_t(S) = \|SA - BS\|_F^2 + t(n - \|S\|_F^2)$$

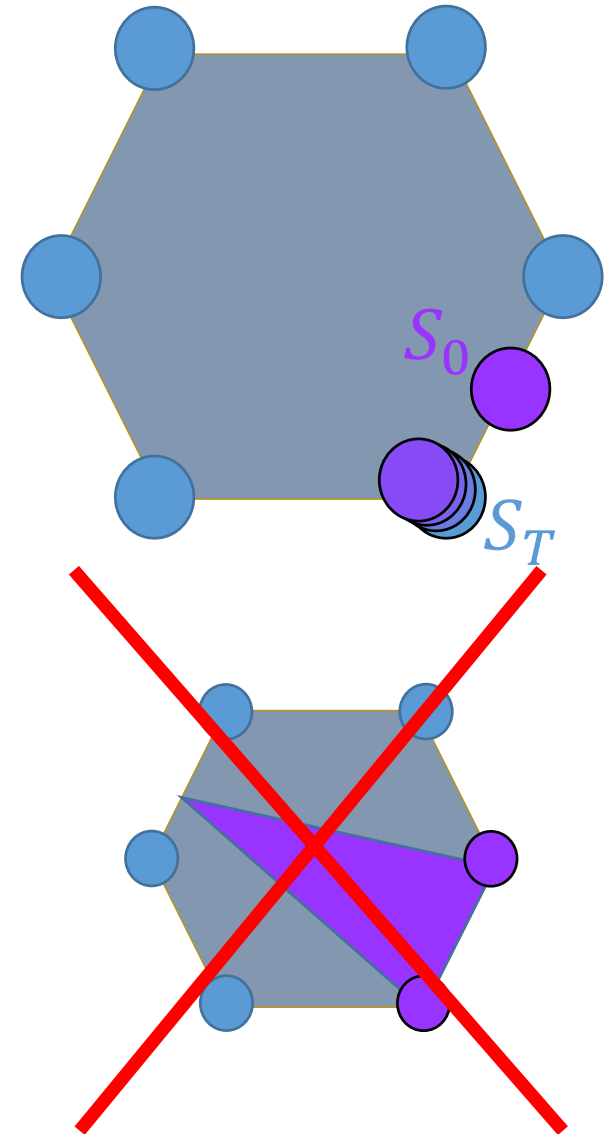
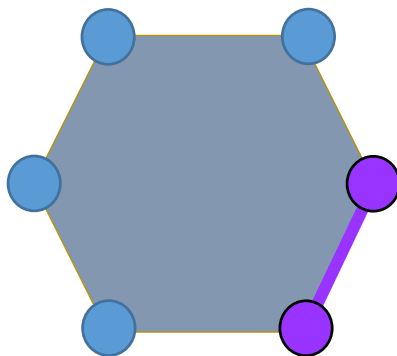
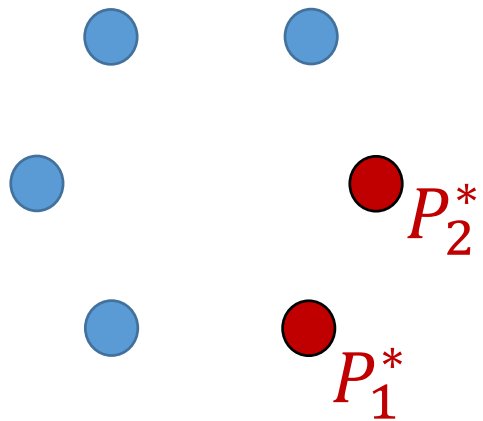
For  $t > 0$ ,  $P_i^*$  are the only global minima!



# Symmetric, no noise

Theorem 3: If  $DS(A,B)$  is convex exact, then  
(under some conditions)

$$S_1, S_2, \dots, S_T = P_i^*$$



# Thank you!

For more details see:

“Exact Recovery with Symmetries for the Doubly-Stochastic Relaxation.”

“DS++: A Flexible, Scalable and Provably Tight Relaxation for Matching Problems.”

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