

# Equal-quantiles rules in resource allocation with uncertain needs

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# Motivation

- Pre-committed division with uncertain needs
- Examples: allocation of public service; division of rescue forces/medical supplies; capacity allocation in a network
- Ex post reallocation may not be possible
- Departure from the literature: waste vs deficit

# The model

- $\mathcal{N}$ : the set of all finite subsets of  $\mathbb{N}$
- $I \in \mathcal{N}$ : a finite population of agents
- $F_i$ : a probability measure on  $\mathbb{R}_+$  with convex and compact support
- $T \in \mathbb{R}_+$ : total endowment
- A problem:  $(F, T) \in \mathcal{P}^I, I \in \mathcal{N}$
- An allocation:  $t \in \mathbb{R}_+^I$  s.t.  $\sum t_i \leq T$  and for each  $i, t_i \in [0, \max \text{supp} F_i]$
- A rule  $r : \bigcup_I \mathcal{P}^I \rightarrow \bigcup_I \mathbb{R}_+^I : r(F, T) = t$

## Cost of an assignment to a single agent $i$

Suppose that  $u_i > u_0 > 0$ . Agent 0 generates deterministic welfare and is outside the model.

Utility maximization  $\iff$  cost minimization

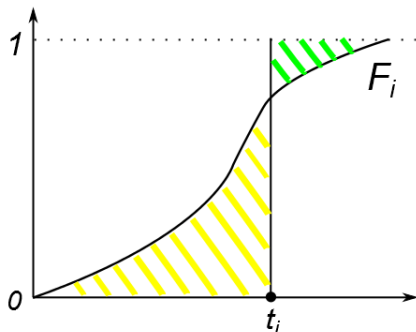
$$\begin{aligned} & \int_{x_i < t_i} u_i x_i + u_0(T - x_i) - [u_i x_i + u_0(T - t_i)] dF_i(x_i) \\ & + \int_{x_i > t_i} u_i x_i + u_0(T - x_i) - [u_i t_i + u_0(T - t_i)] dF_i(x_i) \\ & = \int_{x_i < t_i} u_0(t_i - x_i) dF_i(x_i) \\ & + \int_{x_i > t_i} (u_i - u_0)(x_i - t_i) dF_i(x_i) \\ & = c^w \cdot ew(F_i, t_i) + c_i^d \cdot ed(F_i, t_i) \end{aligned}$$

## Optimal assignment to a single agent $i$

$$\min_{t_i} c^w \cdot ew(F_i, t_i) + c_i^d \cdot ed(F_i, t_i)$$

Unconstrained solution:  $t_i = F_i^{-1}\left(\frac{c_i^d}{c^w + c_i^d}\right)$

Constrained solution:  $t_i = T$



## Discussion 1: Resource may not be exhausted.

Given  $i \in \mathbb{N}$  with  $F_i$ , it could be that

$$t_{F_i}^* := \sup_{T \in \mathbb{R}_+} r_i(F_i, T) < \max \text{supp} F_i.$$

Given  $I \in \mathcal{N}$  with  $F$ , it could be that for each  $i \in I$ ,

$$\sup_{T \in \mathbb{R}_+} r_i(F, T) < \max \text{supp} F_i.$$

Recall: an allocation  $t \in \mathbb{R}_+^I$  is s.t.  $\sum t_i \leq T$  and for each  $i$ ,  $t_i \in [0, \max \text{supp} F_i]$ .

## Discussion 1: Maximum assignment

**Our contribution:** We find that some existing axioms, when extended from deterministic to the uncertain context, imply that for each  $I \in \mathcal{N}$  and each  $(F, T) \in \mathcal{P}^I$ ,

$$T \leq \sum_{i \in I} t_{F_i}^* \Rightarrow \sum_{i \in I} r_i(F, T) = T;$$

$$T > \sum_{i \in I} t_{F_i}^* \Rightarrow \text{for each } i, r_i(F, T) = t_{F_i}^*.$$

Note: The maximum assignment of an agent does not depend on the number of other agents and their claims.

## Discussion 2: Newsvendor problem — Similarity

Each unit of a perishable product can be purchased at price  $c$  and sold at price  $p$  where  $p > c > 0$ .

$$\begin{aligned} & \min_{t_i} c^w \cdot ew(F_i, t_i) + c_i^d \cdot ed(F_i, t_i) \\ &= \min_{t_i} c \cdot ew(F_i, t_i) + (p - c) \cdot ed(F_i, t_i) \end{aligned}$$

Unconstrained solution:

$$t_i = F_i^{-1}\left(\frac{c_i^d}{c^w + c_i^d}\right) = F_i^{-1}\left(\frac{p - c}{p}\right)$$

Critical fractile formula (operations management)

Littlewood's rule (revenue management)



## Discussion 2: Newsvendor problem — Difference

Unlimited resource vs limited resource (multiple agents)

Maximize a utility function vs social choice function

Profit vs social welfare (single agent/multiple agents)

**Our contribution:** Axiomatize a family of division rules selecting allocations according to

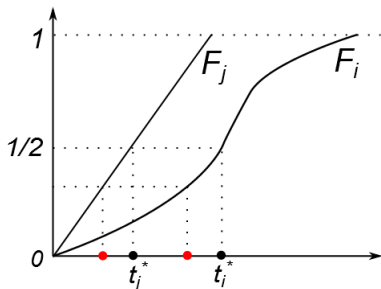
$$\min_t \sum_{i \in I} [c^w \cdot ew(F_i, t_i) + c^d \cdot ed(F_i, t_i)].$$

## Equal-quantile rules

An equal-quantile rule associated with  $\lambda \in (0, 1]$  selects for each  $I \in \mathcal{N}$  and each  $(F, T) \in \mathcal{P}^I$  the allocation that solves

$$\min_t \sum_{i \in I} [c^w \cdot ew(F_i, t_i) + c^d \cdot ed(F_i, t_i)],$$

where  $c^w, c^d > 0$  are such that  $\lambda = \frac{c^d}{c^w + c^d}$ .



## Axioms: *Continuity*

For each  $I \in \mathcal{N}$ ,  $r$  is continuous on  $P^I$ .

## Axioms: *Strict Ranking*

For each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each pair  $i, j \in I$ , if  $F_i$  strictly first-order stochastically dominates  $F_j$ , then  $r_i(F, T) > r_j(F, T)$ .

## Axioms: *Ranking*

For each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each pair  $i, j \in I$ , if  $F_i$  first-order stochastically dominates  $F_j$ , then  $r_i(F, T) \geq r_j(F, T)$ .

*Strict ranking* and *continuity* imply *ranking*.

*Ranking* implies *equal treatment of equals*: For each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each pair  $i, j \in I$ , if  $F_i = F_j$ , then  $r_i(F, T) = r_j(F, T)$ .

## Axioms: *Consistency*

For each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , each  $J \subseteq I$ , and each  $i \in J$ ,

$$r_i(F, T) = r_i(F_J, \sum_{j \in J} r_j(F, T)) \text{ and}$$

$$r_i(F, T) = r_i(F_J, T - \sum_{j \in I \setminus J} r_j(F, T)),$$

where  $F_J$  is the restriction of  $F$  onto  $J$ .

# Maximum assignment result

For each  $i \in \mathbb{N}$  with  $F_i$ , recall  $t_{F_i}^* := \sup_{T \in \mathbb{R}_+} r_i(F_i, T)$ .

## Theorem

*If a rule  $r$  satisfies consistency and continuity, then for each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each  $i \in I$ ,  $r_i(F, T) \leq t_{F_i}^*$ , and if  $T \leq \sum t_{F_i}^*$ ,  $\sum r_i(F, T) = T$ , and thus for each  $i \in I$ ,  $r_i(F, \sum t_{F_i}^*) = t_{F_i}^*$ .*

*If a rule  $r$  satisfies consistency, continuity, and equal treatment of equals, then for each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each  $i \in I$ , if  $T > \sum t_{F_i}^*$ , then  $r_i(F, T) = t_{F_i}^*$ .*

## Axioms: *Ordinality*

Independence of a (non-linear) transformation of the problem due to a common shock.

For each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , each  $\phi : \mathbb{R}_+ \rightarrow \mathbb{R}_+$  that is strictly increasing and continuous, and each  $i \in I$ ,

$$r_i(F^\phi, \sum \phi(r_j(F, T))) = \phi(r_i(F, T)),$$

where for each  $j \in I$  and each  $x_j \in \mathbb{R}_+$ ,  $F_j^\phi(\phi(x_j)) = F_j(x_j)$ .

D'Aspremont and Gevers (1977), Sprumont (1998), Chambers (2007)



## Implication 1: *Scale invariance*

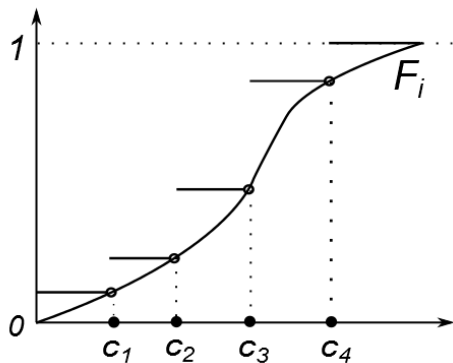
Independence of a uniform rescaling of the problem

Let  $k > 0$  and for each  $x_j \in \mathbb{R}_+$ ,  $\phi^k(x_j) := kx_j$ .

## Implication 2: Coarse ETE

Agents with “coarsely” equal needs receive “coarsely” equal awards.

For each coarse transformation  $\varphi : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ , each  $I \in \mathcal{N}$ , each  $(F, T) \in \mathcal{P}^I$ , and each pair  $i, j \in I$ ,  $F_i^\varphi = F_j^\varphi \Rightarrow \varphi(r_i(F, T)) = \varphi(r_j(F, T))$ .



# Characterization Result

## Theorem

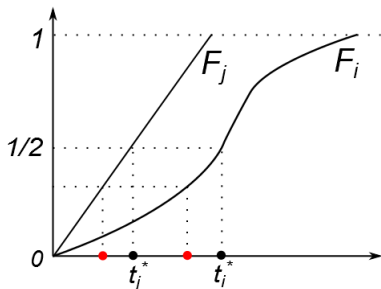
*A rule satisfies continuity, strict ranking, consistency, and ordinality if and only if it is an equal-quantile rule.*

## Equal-quantile rules

An equal-quantile rule associated with  $\lambda \in (0, 1]$  selects for each  $I \in \mathcal{N}$  and each  $(F, T) \in \mathcal{P}^I$  the allocation that solves

$$\min_t \sum_{i \in I} [c^w \cdot ew(F_i, t_i) + c^d \cdot ed(F_i, t_i)],$$

where  $c^w, c^d > 0$  are such that  $\lambda = \frac{c^d}{c^w + c^d}$ .



# Literature

Deterministic fair division: Moulin (2002), Thomson (2003, 2015)

Operations research: Rawls and Turnquist (2010), Wex, Schryen, Feuerriegel, Neumann (2014), etc.

Fair division under uncertainty: Ertemel and Kumar (2018), Xue (2018), Hougaard and Moulin (2018)

# Literature: Fair allocation and welfare economics of risk

Fair allocation: axiomatize division rules (may or may not be rationalizable by some underlying social welfare function)

Welfare economics of risk: axiomatize social welfare functions under risk (Harsanyi (1955), Diamond (1967))

Open question: build a connection between them.